



Multiple Subchannel Sets: An Implementation View

The purpose of this IBM® Redpaper is to demonstrate a pragmatic implementation approach for the adoption of Multiple Subchannel Sets (MSS) in a System z[™] environment.

MSS provides relief for I/O device configurations in large System z10™ and System z9™ environments. It also increases Parallel Access Volumes (PAVs) connectivity.

The following IBM protocols support MSS:

- ► Enterprise Systems Connection (ESCON®)
- ► Fibre Connection (FICON®)
- ► z/OS®

The IBM System Storage™ and the System z™ support PAVs.

MSS description

The following servers support MSS functionality:

- ▶ System z10
- System z9

Note: Do not confuse MSS with multiple Channel Subsystems (CSS). In most cases a subchannel represents an addressable device. For example, a disk control unit with 30 drives uses 30 subchannels. An addressable device is associated with a device number.

Subchannel numbers

Subchannel numbers (including their implied path information to a device) are limited to four hexadecimal digits by hardware and software architectures. These four hexadecimal digits provide 64 K addresses, also known as a set.

IBM has 256 reserved subchannels. That leaves 63.75 K subchannels for general use with the System z10 and System z9 servers.

PAV has made this limitation of subchannels a challenge for larger installations. A single disk drive (with PAV) often consumes at least four subchannels.

Removing these constraints is difficult because four hexadecimal digits for subchannels (and device numbers corresponding to subchannels) are used in a number of places. Simply expanding the field would break too many programs.

The solution gives you the ability to have sets of subchannels (addresses), with a current implementation of two sets. Each set provides 64 K addresses.

- Subchannel set 0 (SS0)
 - Reserves subchannels for IBM use. Although, the number of reserved subchannels
 (256) on the System z10 and System z9 is less than on z990, and z890 servers (1024)
- ► Subchannel set 1 (SS1)
 - Provides a full range of 64 K subchannels on System z10 and System z9 servers.

Note: Each CSS has its own SS0 and SS1.

Figure 1 on page 3 shows the System z server with MSS

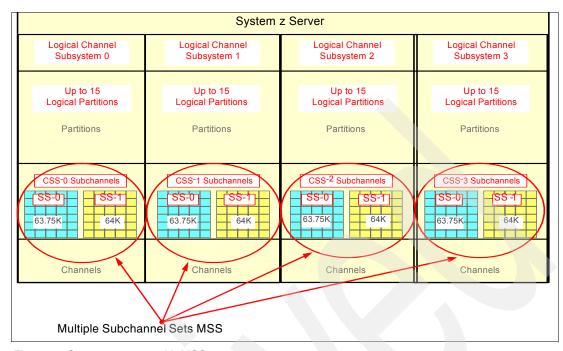


Figure 1 System z server with MSS

MSS-implementation considerations

When you implement MSS, you must understand and have access to the following information:

- ► Hardware requirements
- Software requirements
- ► MSS architecture
- ► Device number alignment
- Migration tools
- Disk subsystem considerations
- Numbering schemes
- PAV implementation

Each implementation consideration is described in detail in the remaining sections of this publication.

Hardware requirements

Only System z10 and System z9 servers support MSS.

Software requirements

MSS requires z/OS 1.7 and above.

Note: Logical partitions (LPARS) with earlier versions of z/OS do not see devices in SS1.

MSS architecture

Correspondence between addresses in the two sets is not required. You can have device number 8000 in SS0, and device number 8000 in SS1.

They might refer to completely separate devices. We know that the device number in SS1 must be an alias for z/OS, but that is all we know from the device number. Likewise, device number 1234 (SS0), and device number 4321 (SS1) might be the base, and an alias for the same device.

An additional high order digit (either a 0 or a 1) is logically added to existing device numbers (e.g. 08000 for device 8000 in SS0 and 18000 for device 8000 in SS1).

This enhancement is not delivered by system code. Remember, there is still an architectural requirement of four-digit addresses (device numbers, subchannels). However, some messages contain subchannel set numbers. You can mentally use the message as a high order digit for device numbers.

Attention: A device address must be unique within a DASD controller. It is made unique by the control unit image (CUADD parameter in the IOCP) and the unit address. This cannot be duplicated across multiple subchannel sets within the same DASD controller.

Device number alignment

Typically, base devices (3390B) are defined from the start of the range (00 forwards). Alias devices are defined from the last device in the range backwards. If the range is filled, the device fast forwards backwards.

Note: Base and alias device addresses on a single control unit must be unique. They cannot be duplicated across subchannel sets.

Plan your device ranges strategically. Devices should be predefined in SS1. This is established using the MAXDEV setting on the channel subsystem definition, which correlates to the RESOURCE statement in the Input/Output Configuration Program (IOCP). This pre-allocates the Hardware System Area (HSA). As z10 servers have a fixed HSA and are always configured for the maximum capability, no preplanning of the MAXDEV setting is required. A power on reset is required to change this setting on System z9 servers.

Migration tools

You can use dynamic reconfiguration to relocate device ranges (UCBs) between SS0 and SS1.

For consolidation of device ranges involving migration of data, various tools are available, such as:

- ▶ DFSMSDSS
 - Migrates data at either a volume level (DEVICE) or logical level (DATASET)
- ► Transparent Data Migration Facility (TDMFTM)
 - Moves data at the block and volume level without disruption

- ► Logical Data Migration Facility (LDMF)
 - Migrates datasets at the volume level
 - Typically used to consolidate smaller volumes

Disk subsystem considerations

We highly recommend that you contact your IBM System Storage representative prior to choosing a numbering scheme for your storage subsystem. Alternatively, contact your vendor for guidance of non-IBM storage subsystems.

Numbering schemes

When choosing a number scheme, you will not find just one correct approach for all circumstances. The approach you choose will vary depending on the customer requirements.

In this section we are going to demonstrate different options for you.

The numbering schemes are the following:

- Scheme 1- Compacting or rolling reuse of addresses
 - Device numbers available to SS0 and SS1 are used on an "as available" basis.
- ► Scheme 2 Direct access storage device (DASD) versus the rest
 - The device numbers used by alias volumes (3390A) are assigned to SS1, then reused in SS0 for non-DASDs.
- Scheme 3 Device Pairing
 - The first half of an LSS range (128 addresses) is allocated to the base (3390B) range.
 The same device range is reused for the next consecutive assignment of devices within SS1.
- Scheme 4 Filling the ranges
 - The device addresses in SS1 are amalgamated to complete ranges, then reused accordingly in SS0.

Scheme 1 - Compacting or rolling reuse of addresses

The numbering scheme duplicates ranges across SS0 and SS1 by using the unit address (UA) field for 3390A in SS1.

- ► 5000,72 UNITADD=00 in SS0
- ▶ 5000,184 UNITADD=48 in SS1

Compaction is achieved in both SS0 and SS1 by using the follow on numbers in each SS (5048 is next in SS0, 50B7 is next in SS1). This achieves maximum reuse of device numbers. However, it should be noted that this does look confusing as the device addresses in SS1 are allocated unit addresses that do not match the numbering scheme (Example 1).

Example 1 IOCP use of the UNITADD for device range compaction

```
IODEVICE ADDRESS=(50B8,184), UNITADD=48, CUNUMBR=(5000), STADET=Y*, DESC='IBM 2107 DASD', UNIT=3390A, SCHSET=1
```

Scheme 1 - Hardware configuration dialog example

Figure 2 on page 6 shows range 1000-1047 defined as 3390B devices in CNTLUNIT 1000.

```
Device / Processor Definition

Row 1 of 1

Command ===> _______ Scroll ===> CSR

Select processors to change device/processor definitions, then press Enter.

Device number . . : 1000 Number of devices . : 72

Device type . . . : 3390B

Preferred Device Candidate List

/ Proc.CSSID SS+ UA+ Time-Out STADET CHPID + Explicit Null

_ MINE.0 _ _ No Yes _ No ____
```

Figure 2 CNTLUNIT 1000: 3390B definitions for UA xx00 - xx47

The remainder of this range (xx48-xxFF) is defined as 3390A devices, but with the same device address (1000). Figure 3 on page 6 shows how this is achieved using the UA field of the device definition.

Note: Notice the alignment to SS1. This is the default for alias devices attached to a System z10 or System z9 within the hardware configuration dialog (HCD).

Figure 3 CNTLUNIT 1000: 3390A definitions for UA xx48 - xxFF

Figure 4 on page 7 shows the CNTLUNIT 1000 I/O device list.

Goto Filter Backup	Query Help		
Command ===>	I/O Device List	Row 1 of 2 More: > Scroll ===> CSR	
Select one or more devic	ces, then press Enter. To add,	use F11.	
Control unit number : 1000 Control unit type . : 2107			
Device	#Control Unit	Numbers +	
/ Number Type +	CSS OS 1 2 3 4 5	5 6 7 8	
1000,72 3390B	1 1000		
_ 1000,184 3390A	1 1000		

Figure 4 ICNTLUNIT 1000 /O device list

The next CNTLUNIT is numbered to align with the next available device address in SS0 (1048). Figure 5 on page 7 shows the next available devices defined as bases (3390B) to SS0.

Note: We specify a UA of 00. This may look strange as we define devices 1048-8F, but you do not want to force a UA of 00 for 1048, 01 for 1049, and so on.

Device / Processor Definition			
	Row 1 of 1		
Command ===>	Scroll ===> CSR		
Select processors to change device/processor definitions, then press Enter. Device number : 1048			
	Preferred Device Candidate List		
/ Proc.CSSID SS+ UA+ Time-Out STADET	CHPID + Explicit Null		
_ MINE.O _ 00 No Yes	No		

Figure 5 CNTLUNIT 1048: 3390B definitions for UA xx00 - xx47

The remainder of the range in CNTLUNIT 1048 is filled with the alias devices in SS1 (UA xx48-xxFF). Figure 6 on page 8 shows how once again we generate a mismatch between the device number and the associated UA (10B8 with UA of 48). Remember, this is perfectly valid.

Device / Processor Definition				
				Row 1 of 1
Command ===>				Scroll ===> CSR
Select processors to ch	ange devic	e/proces	sor definit	ions, then press
Enter.				
Device number : 10	Device number : 10B8 Number of devices . : 184			
Device type : 3390A				
			D C 1	
/ Proc.CSSID SS+ UA+	Timo Out	CTADET		Device Candidate List Explicit Null
MINE.0 1 48		Yes	CULID +	Explicit Null No
_ 111112.0 1 40	110	103	_	

Figure 6 CNTLUNIT 1048: 3390A definitions for UA xx48 - xxFF

Figure 7 on page 8 shows the device list for CNTLUNIT 1048.

```
Goto Filter Backup Query Help
                         I/O Device List
                                                 Row 1 of 2 More:
Command ===>
                                                 Scroll ===> CSR
Select one or more devices, then press Enter. To add, use F11.
Control unit number : 1048
                             Control unit type . : 2107
 ------Device----- --#--- Control Unit Numbers + ------
                    CSS OS 1--- 2--- 3--- 4--- 5--- 6--- 7--- 8---
/ Number Type +
                      1
 1048,72 3390B
                            1048
 10B8,184 3390A
                            1048
```

Figure 7 CNTLUNIT 1048 I/O device list

Scheme 1 - IOCP statements example

Example 2 shows the IOCP statements demonstrating scheme 1 with device range compaction.

Example 2 IOCP statements with device range compaction

```
DESC='IBM 2107 DASD',UNIT=3390B
IODEVICE ADDRESS=(10B8,184),UNITADD=48,CUNUMBR=(1048),STADET=Y*
,DESC='IBM 2107 DASD',UNIT=3390A,SCHSET=1
Ranges aligned to each subchannel set at this point:
SSO: 1000-1047, 1048-108F
SS1: 1000-1087, 1088-116F
```

Scheme 2 - DASD versus the rest

Scheme 2 reuses the freed device numbers in SS0 for all devices other than DASD.

For example, a control unit has:

- ▶ 1000,72 for the base devices (UA=00..47) and
- ▶ 1048,184 for the alias devices in SS1 (UA=48..FF)

The device numbers now allocated to SS1 (1048..10FF) are used to defined device types other than DASD.

Scheme 2 - IOCP statements example

Example 3 shows IOCP statements to demonstrate scheme 2 using device numbers that are used in SS1 for non DASD devices.

Example 3 IOCP statements with reuse of device numbers for non-DASD control units

```
CNTLUNIT CUNUMBR=1000, PATH=((CSS(0),04,08,2C,31,33,37,6F,75)),*

UNITADD=((00,256)), CUADD=0, UNIT=2107

IODEVICE ADDRESS=(1000,72), UNITADD=00, CUNUMBR=(1000), STADET=Y,*

DESC='IBM 2107 DASD', UNIT=3390B

IODEVICE ADDRESS=(1048,184), UNITADD=48, CUNUMBR=(1000), STADET=Y*

,DESC='IBM 2107 DASD', UNIT=3390A, SCHSET=1

Ranges aligned to each subchannel set at this point:
SSO: 1000-1047
SS1: 1048-10FF

Devices 1048-10FF are now available in SSO for use by non-DASD devices.
```

Scheme 3 - Device pairing

When you pair devices, consider the split of base and alias devices. The first half of an LSS range (128 devices) is allocated to the initial base range. You then use this device range again for the next consecutive assignment of devices within SS1.

This scheme only works when no more than 128 alias devices are defined for a range of base devices. Devices are aligned to the UA with this scheme.

Table 1 on page 9 shows how this approach facilitates double the addressable devices.

Table 1 Paired device ranges

RANGES	CNTLUNITS	USED in SS0	USED in SS1
8000-80FF	8000	8000-807F	8080-80FF
	8080	8080-80FF	8000-807F

RANGES	CNTLUNITS	USED in SS0	USED in SS1
8100-81FF	8100	8100-817F	8180-81FF
	8180	8180-81FF	8100-817F

Note: When paired, the UA field for each device will always start at 00 for base addresses and 80 for alias addresses. This shows a difference between the device address number and UA when the range in SS0 starts on the xx80 boundary.

The device range split (base versus alias) adopted in pairing is very conservative in a FICON implementation.

Scheme 3 - HCD example

This section reflects the definition of control units to establish paired device ranges.

CNTLUNIT 8000 definition

Figure 8 on page 10 shows that CNTLUNIT 8000 is defined with the following devices to fill the range:

- ▶ 8000 807F as 3390B devices in SS0
- 8080 80FF as 3390A devices in SS1

Figure 8 CNTLUNIT 8000 I/O device list

CNTLUNIT 8080 definition

Figure 9 on page 11 shows this device range used in CNTLUNIT 8080.

- 8080 80FF as 3390B devices in SS0
- ▶ 8000 807F as 3390A devices in SS1

```
| Solution | Command | Com
```

Figure 9 CNTLUNIT 8080 I/O device list

Scheme 3 - IOCP statements example

Example 4 shows IOCP statements to demonstrate scheme 3 with paired device numbers.

Example 4 IOCP statements with paired device ranges

```
CNTLUNIT CUNUMBR=8000, PATH=((CSS(0), 04, 08, 2C, 31, 33, 37, 6F, 75)), *
               UNITADD=((00,256)),CUADD=0,UNIT=2107
        IODEVICE ADDRESS=(8000,128), UNITADD=00, CUNUMBR=(8000), STADET=Y,*
               DESC='IBM 2107 DASD', UNIT=3390B
         IODEVICE ADDRESS=(8080,128), UNITADD=80, CUNUMBR=(8000), STADET=Y, *
               DESC='IBM 2107 DASD', UNIT=3390A, SCHSET=1
Ranges aligned to each subchannel set at this point
SS0: 8000-807F
SS1: 8080-80FF
CNTLUNIT CUNUMBR=8080, PATH=((CSS(0), 04, 08, 2C, 31, 33, 37, 6F, 75)),*
               UNITADD=((00,256)),CUADD=1,UNIT=2107
        IODEVICE ADDRESS=(8080,128), UNITADD=00, CUNUMBR=(8080), STADET=Y,*
               DESC='IBM 2107 DASD', UNIT=3390B
         IODEVICE ADDRESS=(8000,128), UNITADD=80, CUNUMBR=(8080), STADET=Y,*
               DESC='IBM 2107 DASD', UNIT=3390A, SCHSET=1
Ranges aligned to each subchannel set at this point
SSO: 8000-807F, 8080-80FF
SS1: 8080-80FF, 8000-807F
```

Scheme 4 – Filling the ranges

Scheme 4 preserves the ranges in SS0 and aligns devices to a UA. This scheme is driven by the desired BASE to Alias split that is calculated in advance (potentially via the PAV analysis tool). Basic principles used here are as follows:

- Establish BASE to ALIAS ratio first. Align on to boundaries of 16 devices
- Stick to this alignment for all device ranges
- ► All ranges are filled with 256 devices (00-FF)
- ► Use up all device ranges (a range being an LSS of 256 devices 00-FF)
- Define alias devices to SS1
- ► Reuse devices used in SS1 as bases in SS0
- Reuse devices allocated in SS0 in SS1

Example: 192 Bases to 64 Aliases.

There are 256 available ranges (256 x 256 = 65536 or FFFF). The 64 aliases in SS1 free up 256 x 64 = 16384 devices in SS0 for the new ranges.

16384 / 192 = *85 additional ranges*

Table 2 on page 12 shows how we use this example to create a table showing the available ranges.

Table 2 Scheme 4: Mapping the ranges

Range	CNTLUNIT	Used in SS0	Free in Range SS0	Used in SS1	Free in Range SS1
1000-10FF	1000	1000-10BF	10C0-10FF	10C0-10FF	1000-10BF
1100-11FF	1100	1100-11BF	11C0-11FF	11C0-11FF	1100-11BF
1200-12FF	1200	1200-12BF	12C0-12FF	12C0-12FF	1200-12BF

In this scheme you can see numbers in SS0 that are available for reuse. Remember, there is no allegiance between the device number and the UA. So, in an LSS, we can use the "Spare" device numbers accordingly. Figure 10 on page 12 shows the CNTLUNIT 1000 I/O device list.

Scheme 4 - HCD example

Goto Filter Backup	Query Help		
Command ===>	I/O Device List	Row 1 of 2 More: Scroll ===> CSR	>
Select one or more dev	ices, then press Enter. To a	add, use F11.	
Control unit number :	1000 Control unit type	. : 2107	
Device / Number Type + _ 1000,192 3390B _ 10C0,64 3390A	#Control Un CSS OS 1 2 3 4 1 1000		

Figure 10 CNTLUNIT 1000 I/O device list

With all 256 ranges allocated, you can define unused devices in SS0 to additional CNTLUNITs.

Figure 11 on page 13 shows the reused device ranges in CNTLUNIT 1001. Ranges 10C0-10FF, 11C0-11FF and 12C0-12FF are defined as base addresses in SS0. Range 1000-103F is defined as alias addresses in SS1.

Figure 11 I/O device list for CNTLUNIT 10

Scheme 4 - IOCP statements example

Example 5 shows CNTLUNITs 1000, 1100 and 1200 defined as normal. Similar definitions are made for CNTLUNITs 1100 and 1200.

Example 5 Scheme 4: IOCP statements for CNTLUNIT 1000

You can use the free devices SS0 and SS1 to make complete ranges of 256 devices.

Parallel access volumes

Technology is always changing. This section discusses the different implementations of PAV and how these changes impact the number of aliases we align to our base volumes.

Specific functional detail of PAV is in the following techdoc:

http://www-03.ibm.com/support/techdocs/atsmastr.nsf/WebIndex/TD100311

Dynamic versus static PAVs

The initial implementation of PAV was static PAV. With static PAV, specific alias devices are assigned to bases within the logical subsystem (LSS). You cannot re-assign these aliases to any other base device.

With dynamic PAV, alias devices are unassigned and re-assigned to different base devices. Aliases are assigned on an "as needed" basis with workload manager reacting to performance goals.

Predicting the split of bases and aliases

It is not trivial to predict the optimum ratio of Alias to Base addresses. If the ratio is too large, the amount of physical volumes is limited. If too small, this impacts the input/output supervisor queue (IOSQ) and response time.

Use these basic rules to achieve base to alias alignment:

- 1. Use as many aliases as you can afford.
- 2. Use Table 3 on page 14 for a conservative recommendation. This is typically used for ESCON implementations.

Table 3 Typical ESCON implementation

Number of cylinders	Number of Aliases for Dynamic PAV	Number of Aliases for Static PAV
1 - 3339	1/3	1
3340 - 6678	2/3	2
6679 - 10,017	1	3
10,018 - 16,695	1 1/3	4
16,696 - 23,373	1 2/3	5
23,374 - 30,051	2	6
30,052 - 40,068	2 1/3	7
40,069 - 50,085	2 2/3	8
50,086 - 60,102	3	9
60,193<	3 1/3	10

- 3. Based on the number of channels per LSS. A factor of six times the number of FICON channels allocated to the logical subsystem (LSS) equals the number of recommended aliases
- 4. Allocate I/O rate * response time * 0.002. For example: 5 ms service time with 4000 I/O per second needs 40 aliases.

HyperPAV

With HyperPAV technology, z/OS uses pools of aliases (by LSS). When an application I/O is requested, if the base volume is busy with another I/O, z/OS selects (and removes) a free alias from the pool.

The I/O begins its journey to the base address through the selected alias. When the I/O is done, the alias device is used for another I/O on the LSS, or is returned to the free alias pool. If too many I/Os are started simultaneously, z/OS queues the I/Os at the LSS level. The queued I/O is done first in/first out (FIFO) within assigned I/O priority.

Table 4 on page 15 shows the functional differences.

Table 4 PAV versus HyperPAV functional comparison

Attribute	PAV	HyperPAV
Alias to base ratio	Complex: ➤ One alias per 9GBs, depending on workload and I/O rates and response time requirements ➤ 192 PAV-aliases/64 PAV-bases with Mod 27 volume sizes	Simple: ► (Peak I/O rate * average response time * 2) (10x reduction in the number of PAV-aliases) ► 20 PAV-aliases/236 PAV-bases (10,000/sec * 1 ms * 2)
Workload Management	Sluggish: ► WLM adjustments every 10 seconds when work is not meeting goals, every minute when goals met ► I/O Priority by device Multi-system overheads: ► Aggregate multi-system measurements across SYSPLEX to make correct decisions	On Demand: Instantaneous response to changing work loads and I/O skews across devices ► I/O priority by LSS. No multi-system overhead: ► No multi-system aggregation of measurements needed
RMF™	Number of PAV-aliases bound to a base per interval	Number of times I/O could not start because a PAV-alias was not available, high water mark for PAV-aliases in use for LSS
Improved Efficiency	 Alias used for specific base PAV-aliases bound to the same base across all operating system images Interlocked Device-state-change interrupt required in order to insure multi-system coordination of dynamic PAV-aliases Limited parallelism 	Alias used for any base, any time on any system: Each operating system image uses PAV-alias for a different base at the same time, multiplier for the amount of effective number of aliases No interlock required: Improved parallelism for overcoming the speed of light penalty for replication at distance
VSCR	PAV-aliases reside in 31 bit storage only	10 x reduction in PAV-aliases UCBs and device related data structures

Aliases are used independently in each sysplex z/OS image. The workload manager (WLM) is not involved in alias movement, so it does not need to collect information to manage HyperPAV aliases.

If each LPAR needs 20 aliases, then PAV uses 60 aliases for 3 LPARS.

With HyperPAV, you can service the same requirement with 20 aliases.

A simple HyperPAV analogy:

- ► 4 x LPARs share a DASD LSS
- ► 20,000 I/Os per second go to the LSS
- ► Average response time is 1 millisecond

This means that on average, there are 20 I/Os outstanding at any given time, five from each of the four LPARs. If all requests are for the same BASE device, you need four Aliases.

A total of 10 – 20 are defined to cover spikes, depending on individual comfort zones.

HyperPAV prerequisites

The HyperPAV capability is available with Version 2 of IBM DS8000 storage subsystems. Software levels of z/OS 1.6 and above is necessary with relevant maintenance from the HyperPAV subset of the preventive service planning (PSP) bucket:

http://www14.software.ibm.com/webapp/set2/psp/srchBroker

PAV analysis tool

Use the PAV analysis tool to gauge a better understanding of your workload profile and achieve a more precise split. This tool displays PAV alias usage in a 3-D graphical representation. The tool is available from the following URL:

http://www.ibm.com/servers/eserver/zseries/zos/unix/bpxa1ty2.html

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