

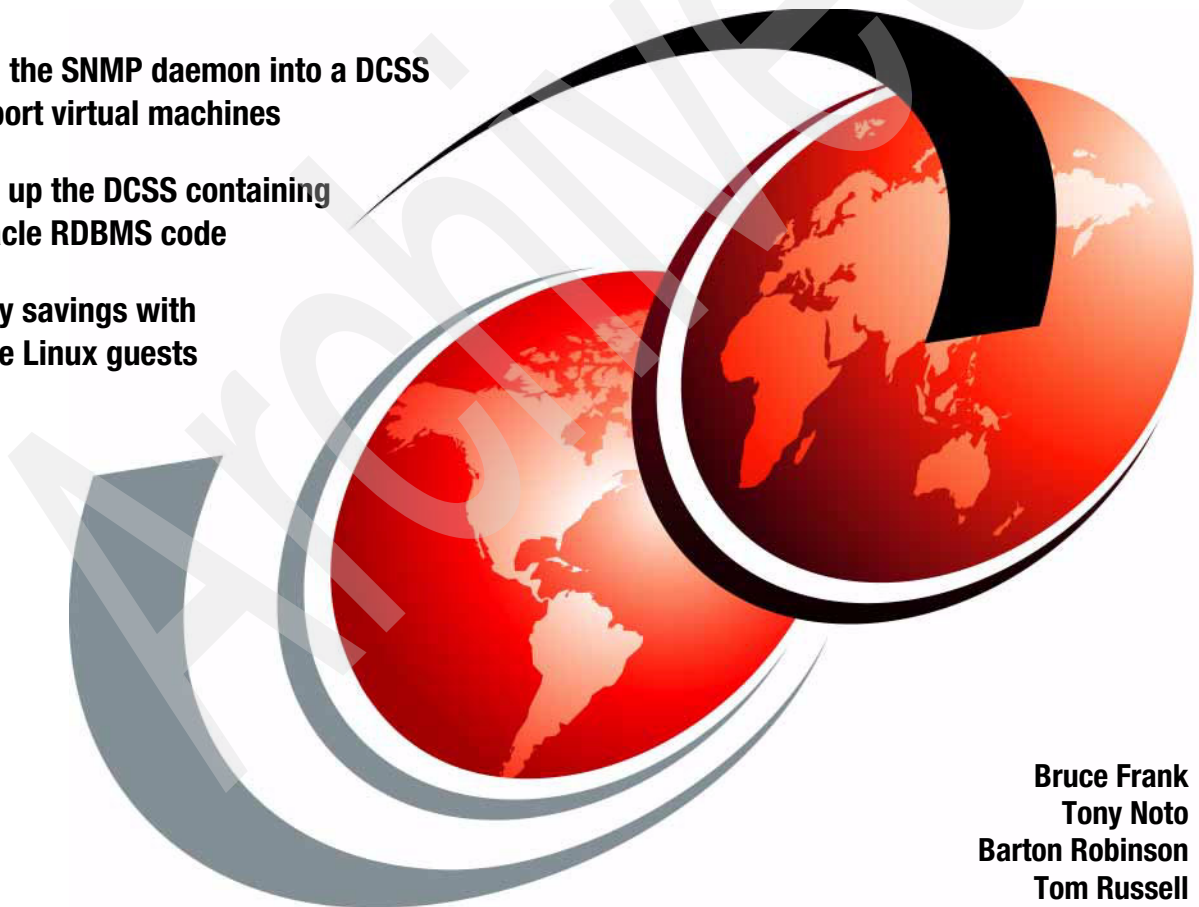
Using Discontiguous Shared Segments and XIP2 Filesystems

With Oracle Database 10g on Linux for IBM System z

Putting the SNMP daemon into a DCSS to support virtual machines

Setting up the DCSS containing the Oracle RDBMS code

Memory savings with multiple Linux guests



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International Technical Support Organization

**Using Discontiguous Shared Segments and XIP2
Filesystems: With Oracle Database 10g on Linux for
IBM System z**

June 2006

Archived

Note: Before using this information and the product it supports, read the information in “Notices” on page v.

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Preface

Linux® on IBM® System z™ offers many advantages to customers who rely on IBM mainframe systems to run their businesses. Linux on IBM System z takes advantage of the qualities of service in the IBM System z hardware and in z/VM® — making it a robust industrial strength Linux. This provides an excellent platform for consolidating Oracle databases that exist in your enterprise.

This IBM Redbook describes our experiences in boosting performance and reducing memory requirements by using discontinuous saved segments (DCSS) and the execute in place (XIP2) filesystem technology.

This book provides an overview of DCSS and XIP2 in Chapter 1, and describes our experiences gained while installing and testing Oracle Database 10g for Linux on IBM System z, namely:

- ▶ Using DCSS and XIP2 with Oracle:
 - Our tests with the Simple Network Management Protocol (SNMP) daemon are described in Chapter 2, “Putting the SNMP daemon into a DCSS” on page 5.
 - Our tests with Oracle are described in Chapter 3, “Using DCSS for Oracle” on page 23.
 - The measurements we made are described in Chapter 4, “DCSS measurements” on page 37.

Interested readers include database consultants, installers, administrators, and system programmers.

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DCSS and XIP2 overview

This chapter provides an overview of DCSS and XIP2:

- ▶ 1.1, “Using DCSS in Linux SLES9” on page 2
- ▶ 1.2, “Using the XIP2 filesystem in Linux SLES9” on page 2
- ▶ 1.3, “DCSS for SNMP” on page 3
- ▶ 1.4, “DCSS for Oracle modules” on page 3
- ▶ 1.5, “Oracle and XIP2 considerations” on page 4

In further chapters, we provide information about using DCSS in an Oracle environment with these topics:

- ▶ Our tests with the SNMP daemon are described in Chapter 2, “Putting the SNMP daemon into a DCSS” on page 5.
- ▶ Our tests with Oracle described in Chapter 3, “Using DCSS for Oracle” on page 23.
- ▶ The measurements we made described in Chapter 4, “DCSS measurements” on page 37.

1.1 Using DCSS in Linux SLES9

The high level objective of this project is to document techniques that use features of z/VM which a customer can implement that can reduce the storage requirement for the second and subsequent Linux servers.

Each Linux instance under z/VM loads its own copies of programs, and manages its own storage. When multiple copies of Linux are running on the same copy of z/VM there is an opportunity to use functions unique to z/VM and System z to share memory pages between the multiple Linux machines.

A Discontiguous Shared Segment (DCSS) is a function of z/VM that allows a segment of memory to be loaded into a virtual machine by command. This segment of memory can be discontiguous to the virtual machine's address space, and a shared copy is loaded at the same address in all virtual machines that load the same DCSS.

A segment is a 1 MB portion of real storage as defined by the ESA/390 architecture. A z/VM saved segment is a range of pages of virtual storage consisting of a number of memory segments where you can hold data or reentrant code (programs). A discontiguous saved segment (DCSS) is a saved segment that can be embedded above the virtual machine's defined storage size.

In a virtual server farm with similar Linux instances there is often a considerable amount of data that is required by all instances. A DCSS enables z/VM to load such data into a designated segment of physical memory and allows all Linux instances that need the data to share this memory segment.

1.2 Using the XIP2 filesystem in Linux SLES9

The execute in place (XIP2) filesystem is provided in SuSE SLES9. This filesystem allows programs residing in a DCSS to be executed directly, sharing the resident pages with other Linux guests.

You can find the latest version of the book *How to use Execute-in-Place Technology with Linux on z/VM*, SC33-8287 at:

http://www-128.ibm.com/developerworks/linux/linux390/october2005_documentation.html

If you are routed to the top-level page, on the left navigation bar, under *Kernel 2.6*, click **October 2005 stream**. On the October 2005 stream page, click **Documentation**.

The book discusses sharing complete directories in a DCSS as well as using a DCSS to share individual files. We will show both here in this redbook. You can also download a tarball containing shell scripts that you can tailor to set up the DCSS environment. We modified these scripts slightly for our environment. The modified versions are listed in Appendix B, “Scripts used to set up DCSS” on page 65.

1.3 DCSS for SNMP

Every Linux guest uses the SNMP daemon. This daemon listens on TCP/IP port 161 and will reply with data on the activity in the Linux guest. Our monitor, ESAMON, sends SNMP requests to each active guest every minute. By putting the SNMP daemon into a DCSS, each Linux guest will use the same copy of the pages where the program is loaded.

Plan

As discussed in Chapter 2, we analyzed the amount of disk space that the SNMP daemon takes, and decided that 8 MB was enough to hold the shared segment containing the SNMP daemon. We also decided to use the technique for sharing individual files as opposed to sharing directories.

The DCSS is executable code. In order to address these pages, Linux must have built page and segment tables to address the memory where the segment will be loaded. These memory management data structures are normally only built for the amount of memory in the virtual machine. Our Linux guests were 256 MB each.

Define

Under z/VM we defined the segment NETSNMP. We reduced the size of the virtual machine by 8 MB and defined the DCSS to be loaded with an end address of 256 MB. The Linux boot parms were changed to ensure that the Linux guest could address storage from zero to 256 MB. See Chapter 2, “Putting the SNMP daemon into a DCSS” on page 5 for the details.

1.4 DCSS for Oracle modules

Plan

After many experiments, we decided that the most effective way to use the DCSS for Oracle is to share an entire directory, and only put into the DCSS the code that we knew would be executed. In our tests we got the best results by putting only two modules (`oracle` and `sqlplus`) in the DCSS.

The standard scripts for implementing DCSS as supplied from developerWorks® assume that the code loaded in the DCSS is in the root file system. Because \$ORACLE_HOME is not in the root file system, we modified the techniques slightly. Chapter 3, “Using DCSS for Oracle” on page 23 describes the process we used for defining and implementing this DCSS.

1.5 Oracle and XIP2 considerations

The execute in place (XIP2) filesystem allows Linux servers to share programs, executing the code directly out of the DCSS. The objective is that with many Linux servers sharing storage, programs only need to be resident once and should allow for significant storage savings.

There are trade-offs that must be understood before implementing XIP2. In order to access the DCSS, the Linux kernel must be configured with enough virtual and real page management structures to address the DCSS pages. These additional control blocks can sometimes cost more than the savings generated by sharing the DCSS. This is discussed in detail in Chapter 4, “DCSS measurements” on page 37.

ESALPS Version 4.1, the Linux Performance Suite, was provided by Velocity Software for this analysis. More information about ESALPS can be found in Appendix A, “ESALPS” on page 61.

Putting the SNMP daemon into a DCSS

This chapter documents the process that we used to put selected files that support the SNMP daemon in Linux. Because every Linux machine uses SNMP, we put it into a DCSS so that a single copy could support all of the virtual machines.

In this chapter, we discuss:

- ▶ Creating the DCSS segment
- ▶ Attaching the DCSS segment to the Linux guest, creating a file system on the DCSS and putting the programs into it that we want to share
- ▶ Making changes to the Linux guest so that the DCSS is allocated at boot time, and correcting files that are overmounted by the DCSS version.
- ▶ Rebooting the machine to begin using the DCSS copy of the shared routines.

2.1 DCSS creation for the SNMP daemon

The following is the procedure used for creating the “direct file” DCSS.

The DCSS creation requires a virtual machine with Class E privileges. There are other methods of doing this without giving the Linux machine authority class E. For this project, we found it simpler to give the Linux machine class E, and do all of the CP segment definition from the Linux machine.

Information about how to use the execute in place (XIP2) filesystem was received from the developerWorks site for documentation and sample code:

http://www-128.ibm.com/developerworks/linux/linux390/october2005_documentation.html

A tarball of the sample scripts that we used can be downloaded from here.

2.2 Sizing the DCSS

To make effective use of the DCSS, you should ensure that it is sized correctly. The following is the procedure that we used to build a DCSS that was the minimum size to hold the number of files that we knew would be used. The shell scripts that we used, as well as the methodology, was downloaded from the developerWorks Web site referred to in Figure . The scripts were put in the home directory of the user doing the work.

To determine the size of the processes, run the **copylibs.sh** script that was downloaded from the developerWorks Web site. This execution will copy all snmp programs from the current \$PATH to a directory /dcss. For example, the following is used to copy the SNMP daemon programs.

Example 2-1 Sample execution of copylibs.sh

```
linux9:~ # xip-howto-scripts/copylibs.sh -f snmpd -d /dcss
option f is set with snmpd
snmpd
option d is set with /dcss
/dcss
/usr/local/lib/libnetsnmpmibs.so.5
/usr/local/lib/libnetsnmpmibs.so.5 is a link
/usr/local/lib/libnetsnmpmibs.so.5.2.1
/usr/local/lib/libnetsnmpagent.so.5
/usr/local/lib/libnetsnmpagent.so.5 is a link
/usr/local/lib/libnetsnmpagent.so.5.2.1
/usr/local/lib/libnetsnmphelpers.so.5
/usr/local/lib/libnetsnmphelpers.so.5 is a link
```

```
/usr/local/lib/libnetsnmphelpers.so.5.2.1
/usr/local/lib/libnetsnmp.so.5
/usr/local/lib/libnetsnmp.so.5 is a link
/usr/local/lib/libnetsnmp.so.5.2.1
/lib/libdl.so.2
/usr/lib/librpm-4.1.so
/usr/lib/librpmdb-4.1.so
/usr/lib/librpmio-4.1.so
/lib/tls/librt.so.1
/lib/tls/libpthread.so.0
/lib/libz.so.1
/lib/libz.so.1 is a link
/lib/libz.so.1.2.1
/usr/lib/libpopt.so.0
/usr/lib/libpopt.so.0 is a link
/usr/lib/libpopt.so.0.0.0
/usr/lib/libbz2.so.1
/usr/lib/libbz2.so.1 is a link
/usr/lib/libbz2.so.1.0.0
/lib/tls/libm.so.6
/lib/tls/libc.so.6
/lib/ld.so.1
/lib/ld.so.1 is a link
/lib/ld-2.3.3.so
```

Note: Libraries which were loaded with `libdl` will not be copied, be sure that you copy these manually.

The execution of the script in Example 2-1 on page 6 copies all the SNMP daemon files to a new directory `/dcss`. Now, to determine how much space the files take up, issue the `du -sk` command, as shown in Example 2-2. This will show how much space (in kilobytes) is used in the `/dcss` directory.

Example 2-2 Calculating size of /dcss directory

```
linux9:~ # du -sk /dcss
7024    /dcss
```

Each file has some overhead. Adding 4 KB per file provides for this additional overhead. To determine how many files there are in the DCSS, issue the command as shown in Example 2-3:

Example 2-3 Counting the files to go in the DCSS

```
linux9:~ # find /dcss type f -print | wc -l
17
```

The DCSS is allocated in one MB increments so the size will be rounded up to the next whole megabyte. Now with the size of the programs loaded, and the number of files, the DCSS storage requirement can be calculated as demonstrated in Example 2-4.

Example 2-4 Calculating the size of the DCSS

Number of files * 4 KB = 17 * 4 = 68 KB
Size of files + file overhead = 7024 KB + 68 KB = 7092 KB

Rounding 7092 KB up to the next MB results in 8 MB. Thus the DCSS needs to be 8 MB to support the SNMP daemon process used in this example.

2.3 Determining DCSS addresses

As the DCSS in this example is 8 MB, the next step is to determine the starting address of the DCSS. For this example, the virtual machine size is 256 MB. To get the most benefit in this case from the DCSS, the virtual machine will be reduced in size by 8 MB and the DCSS will be placed just above the virtual machine. The DCSS starting address will be:

$256 \text{ MB} - 8 \text{ MB} = 248 \text{ MB}$

For creating the DCSS, the page number must be determined to supply to the **CP command** that creates the DCSS. The starting address of 248 MB must be converted to hex, and this is then converted to the page number. The DCSS start address of 248 MB is address hex "F800000". The end address is hex "FFFFFFF", or 256 MB minus 1. The page address is then the hexadecimal value truncated, so in this case, the starting page is F800, and ending page is FFFF.

For different applications, a different sized DCSS might be more appropriate.

2.4 Tailoring Linux storage size to support the DCSS

The virtual machine should now be reduced by 8 MB, and then Linux needs to be configured to address the DCSS. By default, Linux will build page management control blocks for the amount of real memory that it detects when it boots in the virtual machine. To increase this value, so that the pages to be used by the DCSS will be addressable, the configuration file "/etc/zipl.conf" requires modification. To do this, we use the standard **vi**, as shown in Example 2-5 on page 9.

Example 2-5 Changing the zipl.conf file

```
cd /etc
vi zipl.conf
```

add “**mem=256M**” to the parameter line (shown in Example 2-6):

Example 2-6 zipl.conf file

```
parameters = "dasd=0.0.0205,0.0.0306 root=/dev/dasda1 selinux=0
TERM=dumb elevator=cfq mem=256M"
```

Run **zipl** to update everything. The output from the **zipl** command is shown in Example 2-7.

Example 2-7 Output from zipl command

```
linux9s:/etc # zipl
Using config file '/etc/zipl.conf'
Building bootmap in '/boot/zipl'
Adding IPL section 'ipl' (default)
Preparing boot device: dasda (0205).
Done.
```

You may want to use the `-V` parameter to get details on `zipl`'s actions. The output from the `zipl` command with the `-V` parameter is shown in Example 2-8.

Example 2-8 Output from `zipl -V` command

```
linux9s:/etc # zipl -V
Using config file '/etc/zipl.conf'
Target device information
Device.....: 5e:00
Partition.....: 5e:01
Device name.....: dasda
DASD device number.....: 0205
Type.....: disk partition
Disk layout.....: ECKD/compatible disk layout
Geometry - heads.....: 15
Geometry - sectors.....: 12
Geometry - cylinders.....: 1113
Geometry - start.....: 24
File system block size.....: 4096
Physical block size.....: 4096
Device size in physical blocks...: 200316
Building bootmap in '/boot/zipl'
Adding IPL section 'ipl' (default)
  kernel image.....: /boot/image at 0x10000
  kernel parmline...: 'dasd=0.0.0205,0.0.0306 root=/dev/dasda1 selinux=0
TERM=dumb elevator=cfq mem=256M' at 0x1000
  initial ramdisk...: /boot/initrd at 0x800000
Preparing boot device: dasda (0205).
Syncing disks...
Done.
```

2.5 Creating the DCSS

We chose to create the DCSS from the Linux Virtual Machine. In order to have the authority to create and populate a DCSS, the z/VM directory entry for this Linux machine must define the Linux guest as having Class E privileges. There must also be a **NAMESAVE** statement in the directory entry. These privileges allow the Linux machine to define the DCSS. In Example 2-9 on page 11, we called the DCSS “NETSNMP”. The parameters for the `defseg` command were calculated in 2.3, “Determining DCSS addresses” on page 8.

Example 2-9 Defining the DCSS

```
#CP DEFSEG NETSNMP F800-FFFF SR
HCPNSD440I Saved segment NETSNMP was successfully defined in fileid 0149.
```

Validate the NSS is done with the QUERY command as shown in Example 2-10.

Example 2-10 Saving the DCSS to create the storage image

```
#CP Q NSS MAP
FILE FILENAME FILETYPE MINSIZE BEGPAG ENDPAG TYPE CL #USERS PARMREGS VMGROUP
0096 DEMOMON DCSS N/A 04000 04FFF SN R 00003 N/A N/A
0092 ESAWEB DCSS N/A 03800 038FF SN R 00007 N/A N/A
0097 ZMON DCSS N/A 02800 037FF SN A 00000 N/A N/A
...
0149 NETSNMP DCSS N/A 0F800 0FFFF SR S 00000 N/A N/A
```

Note that after the **saveseg** command (see Example 2-11), the class of the DCSS has changed from "S" for Skeleton to "A" for Active.

Example 2-11 Results of the "saveseg netsnmp" command

```
#CP SAVESEG NETSNMP
HCPNSS440I Saved segment NETSNMP was successfully saved in fileid 0149.
CP Q NSS MAP
FILE FILENAME FILETYPE MINSIZE BEGPAG ENDPAG TYPE CL #USERS PARMREGS VMGROUP
0096 DEMOMON DCSS N/A 04000 04FFF SN R 00003 N/A N/A
0092 ESAWEB DCSS N/A 03800 038FF SN R 00007 N/A N/A
0097 ZMON DCSS N/A 02800 037FF SN A 00000 N/A N/A
...
0149 NETSNMP DCSS N/A 0F800 0FFFF SR A 00000 N/A N/A
```

2.6 Resize your virtual machine and then reboot

Now that you have saved the DCSS, you have to change the size of your virtual machine so that its pages will not overlap with the DCSS pages when the DCSS is loaded. To do this, shutdown your system, resize the virtual machine, and then reboot. Example 2-12 demonstrates the command required to do this.

Example 2-12 Shutdown, change storage size and reboot

```
linux9s:~ # shutdown -h now
Broadcast message from root (pts/0) (Thu Dec 22 10:57:27 2005):
```

```
The system is going down for system halt NOW!
...
The system will be halted immediately.
...
Power down.
HCPGSP2630I The virtual machine is placed in CP mode due to a SIGP stop
and store status from CPU 00.
DEF STOR 248M
STORAGE = 0248M
Storage cleared - system reset.
IPL 205
```

Changing the virtual machine size with the define storage command works immediately, but the change will be lost if you logoff the virtual machine. To make this a permanent change, update the CP directory entry for this virtual machine. The size of the machine is on the USER statement. This will ensure that the virtual machine size will be correct through system restarts.

After your system is back up, ensure the "mem=" parameter is present and correct by using the command, `cmdline`, as shown in Example 2-13.

Example 2-13 Verify zipl changes were done correctly

```
linux9s:~ # cat /proc/cmdline
dasd=0.0.0205,0.0.0306 root=/dev/dasda1 selinux=0 TERM=dumb elevator=cfq
mem=256M BOOT_IMAGE=0
```

2.7 Setup for the XIP2 DCSS

One of the shell scripts downloaded from the developerWorks Web site shown in Figure on page 6 is "xipinit-fw". Copy this script to /sbin and modify it to point to the eventual mount point of the XIP2 file system and to tell it the name of the DCSS (in this case "NETSNMP"). This is demonstrated in Example 2-14.

Example 2-14 Customizing xipinit-fw

```
cp /root/xip-howto-scripts/xipinit-fw.sh /sbin/xipinit-fw
vi /sbin/xipinit-fw
...
#mount point of xipimage
MPXIPIMAGE="/xip"
#name of xipimage
XIPIMAGE="NETSNMP"
```

Create the mount point with the `mkdir` command (see Example 2-15).

Example 2-15 Create the directory for the mount point

```
mkdir /xip
```

Now load the DCSS driver with the `modprobe` command. This driver is supplied in the Linux distribution, and allows you to treat the DCSS as a disk so that you can read and write data to it. We will use this driver to create a filesystem on the DCSS and copy the files we want to share into this file system. The command in Example 2-16 loads the block driver.

Example 2-16 Load the DCSS block mode driver

```
linux9s:~ # modprobe dcssblk
```

In order to define devices to use this driver we need to define some nodes. To do this we need the driver's major number. To determine this number we look at the `/proc/devices` file. Example 2-17 shows us that the `dcssblk` driver is there and that the major number is 252.

Example 2-17 Find the driver's major number

```
linux9s:~ # cat /proc/devices
Character devices:
 1 mem
 4 ttyS
 4 ttysclp
 5 /dev/tty
 5 /dev/console
 5 /dev/ptmx
10 misc
43 ctctty
128 ptm
136 pts
227 ttyTUB
228 fs3270
Block devices:
 1 ramdisk
 7 loop
 8 sd
 9 md
65 sd
66 sd
67 sd
68 sd
69 sd
70 sd
```

```
71 sd
94 dasd
128 sd
129 sd
130 sd
131 sd
132 sd
133 sd
134 sd
135 sd
252 dcssblk
253 device-mapper
254 mdp
```

The next step shown in Example 2-18 is to create the device nodes, if required.

Example 2-18 Verify device nodes and create if necessary

```
linux9s:~ # ls -l /dev/dcss*
/bin/ls: /dev/dcss*: No such file or directory
```

```
linux9s:~ # mknod /dev/dcssblk0 b 252 0
linux9s:~ # mknod /dev/dcssblk1 b 252 1
linux9s:~ # mknod /dev/dcssblk2 b 252 2
linux9s:~ # mknod /dev/dcssblk3 b 252 3
```

```
linux9s:~ # ls -l /dev/dcss*
brw-r--r-- 1 root root 252, 0 Oct 12 17:15 /dev/dcssblk0
brw-r--r-- 1 root root 252, 1 Oct 12 17:16 /dev/dcssblk1
brw-r--r-- 1 root root 252, 2 Oct 12 17:16 /dev/dcssblk2
brw-r--r-- 1 root root 252, 3 Oct 12 17:16 /dev/dcssblk3
```

2.8 Load the DCSS storage with the files to be shared

At this point, we have created a DCSS segment that will be loaded from address 248 MB to 256 MB. The machine is defined to be 248 MB in size and rebooted at this size. The Linux command line has been modified to support addresses up to 256 MB.

The `dcssblk` driver is used to manipulate the DCSS (Example 2-19).

Example 2-19 Use the `dcssblk` driver to load the DCSS into storage

```
linux9s:~ # echo "NETSNMP" > /sys/devices/dcssblk/add
extmem info:segment_load: loaded segment NETSNMP range 0f800000 ..0fffffff type
SR in shared mode
```

```
dcssblk info: Loaded segment NETSNMP, size = 8388608 Byte, capacity = 16384 512
byte) sectors
Dec 12 17:06:14 linux9s kernel: extmem info:segment_load: loaded segment
NETSNMP range 0f800000 .. 0fffffff type SR in shared mode
Dec 12 17:06:14 linux9s kernel: dcssblk info: Loaded segment NETSNMP,
size = 838 8608 Byte, capacity = 16384 (512 Byte) sectors
```

We now must get exclusive access to the DCSS while we are writing into it. We again use the dcssblk driver to do this (Example 2-20).

Example 2-20 Get exclusive access to the DCSS

```
linux9s:~ # echo 0 > /sys/devices/dcssblk/NETSNMP/shared
```

Now that we have exclusive access, we can define a filesystem on the DCSS (Example 2-21).

Example 2-21 Make a file system on the DCSS

```
linux9s:~ # mke2fs -b 4096 /dev/dcssblk0
mke2fs 1.36 (05-Feb-2005)
Filesystem label= OS type: Linux
Block size=4096 (log=2)
Fragment size=4096 (log=2)
2048 inodes, 2048 blocks
102 blocks (4.98%) reserved for the super user
First data block=0
1 block group
32768 blocks per group, 32768 fragments per group
2048 inodes per group
Writing inode tables: done
Writing superblocks and filesystem accounting information: done
This filesystem will be automatically checked every 37 mounts or
180 days, whichever comes first. Use tune2fs -c or -i to override.
linux9s:~ #
```

Example 2-22 shows mounting the new file system on /xip and copying the files into it.

Example 2-22 Copy of new file

```
linux9:~ # mount /dev/dcssblk0 /xip
linux9s:/dcss # cp -va /dcss/* /xip
'/dcss/lib' -> '/xip/lib'
'/dcss/lib/libdl.so.2' -> '/xip/lib/libdl.so.2'
'/dcss/lib/tls' -> '/xip/lib/tls'
'/dcss/lib/tls/librt.so.1' -> '/xip/lib/tls/librt.so.1'
'/dcss/lib/tls/libpthread.so.0' -> '/xip/lib/tls/libpthread.so.0'
'/dcss/lib/tls/libm.so.6' -> '/xip/lib/tls/libm.so.6'
'/dcss/lib/tls/libc.so.6' -> '/xip/lib/tls/libc.so.6'
'/dcss/lib/libz.so.1.2.1' -> '/xip/lib/libz.so.1.2.1'
'/dcss/lib/ld-2.3.3.so' -> '/xip/lib/ld-2.3.3.so'
'/dcss/usr' -> '/xip/usr'
'/dcss/usr/local' -> '/xip/usr/local'
'/dcss/usr/local/sbin' -> '/xip/usr/local/sbin'
'/dcss/usr/local/sbin/snmpd' -> '/xip/usr/local/sbin/snmpd'
'/dcss/usr/local/lib' -> '/xip/usr/local/lib'
'/dcss/usr/local/lib/libnetsnmpmibs.so.5.2.1' ->
'/xip/usr/local/lib/libnetsnmpmibs.so.5.2.1'
'/dcss/usr/local/lib/libnetsnmpagent.so.5.2.1' ->
'/xip/usr/local/lib/libnetsnmpagent.so.5.2.1'
'/dcss/usr/local/lib/libnetsnmphelpers.so.5.2.1' ->
'/xip/usr/local/lib/libnetsnmphelpers.so.5.2.1'
'/dcss/usr/local/lib/libnetsnmp.so.5.2.1' ->
'/xip/usr/local/lib/libnetsnmp.so.5.2.1'
'/dcss/usr/lib' -> '/xip/usr/lib'
'/dcss/usr/lib/librpm-4.1.so' -> '/xip/usr/lib/librpm-4.1.so'
'/dcss/usr/lib/librpmdb-4.1.so' -> '/xip/usr/lib/librpmdb-4.1.so'
'/dcss/usr/lib/librpmio-4.1.so' -> '/xip/usr/lib/librpmio-4.1.so'
'/dcss/usr/lib/libpopt.so.0.0.0' -> '/xip/usr/lib/libpopt.so.0.0.0'
'/dcss/usr/lib/libbz2.so.1.0.0' -> '/xip/usr/lib/libbz2.so.1.0.0'
```

Now save the DCSS (Example 2-23). This will take the data in the current storage copy of the DCSS and schedule it for writing out to the z/VM SPOOL. Again we use the `dcssblk` driver to do this.

Example 2-23 Save the segment

```
linux9:~ # echo 1 > /sys/devices/dcssblk/NETSNMP/save
Dec 12 17:06:14 linux9s kernel: dcssblk info: Loaded segment NETSNMP,
size = 838 8608 Byte, capacity = 16384 (512 Byte) sectors
dcssblk info: Segment NETSNMP is currently busy, it will be saved when
it becomes idle...
Dec 12 17:17:41 linux9s kernel: dcssblk info: Segment NETSNMP is
currently busy, it will be saved when it becomes idle...
```

Unmount the file system. This means that no-one is connected to the DCSS any more, so it will be written out (Example 2-24).

Example 2-24 Free the segment so that it will be written

```
umount /xip
HCPNSD440I Saved segment NETSNMP was successfully defined in fileid 0150.
HCPNSS440I Saved segment NETSNMP was successfully saved in fileid 150.
dcssblk info: Segment NETSNMP became idle and is being saved now
Dec 12 17:20:26 linux9s kernel: dcssblk info: Segment NETSNMP became
idle and is being saved now
Finish unloading it:
```

We have now finished using the segment as a disk. We can release it, and then mount it using the new XIP2 file system. Example 2-25 demonstrates how this is done.

Example 2-25 Release the DCSS

```
echo NETSNMP > /sys/devices/dcssblk/remove
linux9s:/dcss # ls -l /sys/devices//dcss*
total 0
drwxr-xr-x  2 root root    0 Dec 12 17:25 .
drwxr-xr-x  7 root root    0 Dec 12 16:56 ..
--w-----  1 root root    0 Dec 12 17:06 add
-rw-r--r--  1 root root 4096 Dec 12 17:06 detach_state
--w-----  1 root root    0 Dec 12 17:25 remove
```

Note that the DCSS has been saved into a new file. You can now remove the original one (Example 2-26) to free the space in SPOOL used by the original DCSS.

Example 2-26 Delete the old segment

```
#cp purge nss 149
```

Now that the DCSS has been written out, and it contains the files we wanted, we can mount it. Example 2-27 demonstrates mounting the DCSS with the XIP2 file system.

Example 2-27 Mount the DCSS with the XIP2 file system

```
linux9s:~ # echo "NETSNMP" > /sys/devices/dcssblk/add  
linux9s:~ # mount -t ext2 -o ro,xip /dev/dcssblk0 /xip
```

On the Virtual Machine console you get the messages that the segment was loaded (Example 2-28).

Example 2-28 Messages from loading segment

```
extmem info:segment_load: loaded segment NETSNMP range 0f800000 ..  
0fffffff type SR in shared mode  
Dec 12 17:28:08 linux9s kernel: extmem info:segment_load: loaded segment  
NETSNMP  
range 0f800000 .. 0fffffff type SR in shared mode
```

Check /proc/mounts to verify that the filesystem is mounted at the /xip mount point (Example 2-29).

Example 2-29 Checking the mount points

```
linux9s:/dcss # cat /proc/mounts  
rootfs / rootfs rw 0 0  
/dev/root / ext3 rw 0 0  
proc /proc proc rw 0 0  
sysfs /sys sysfs rw 0 0  
devpts /dev/pts devpts rw 0 0  
tmpfs /dev/shm tmpfs rw 0 0  
none /xip xip2 ro 0 0
```

These files should now be available for any use. You can verify that they are readable by using the **tar** command to copy them (to /dev/null). Example 2-30 demonstrates how to get to the **tar** command and use it to verify that the files are, indeed, accessible.

Example 2-30 Verify files are accessible

```
linux9s:/dcss # cd /xip
linux9s:/xip # tar cvf /dev/null *
lib/
lib/libdl.so.2
lib/tls/
lib/tls/librt.so.1
lib/tls/libpthread.so.0
lib/tls/libm.so.6
lib/tls/libc.so.6
lib/libz.so.1.2.1
lib/ld-2.3.3.so
lost+found/
usr/
usr/local/
usr/local/sbin/
usr/local/sbin/snmpd
usr/local/lib/
usr/local/lib/libnetsnmpmibs.so.5.2.1
usr/local/lib/libnetsnmpagent.so.5.2.1
usr/local/lib/libnetsnmphelpers.so.5.2.1
usr/local/lib/libnetsnmp.so.5.2.1
usr/lib/
usr/lib/librpm-4.1.so
usr/lib/librpmdb-4.1.so
usr/lib/librpmio-4.1.so
usr/lib/libpopt.so.0.0.0
usr/lib/libbz2.so.1.0.0
```

Now that we know that the DCSS is connected and files in it can be used, we test the "overmounting". This is done so that when any of the programs in this list are fetched, the DCSS copy will be used. The overmounting is done by the "xipinit-fw" shell script that we downloaded from the developerWorks Web site. Example 2-31 demonstrates using the xipinit-fw shell script to test the overmounting.

Example 2-31 Overmount the modules in the DCSS

```
umount /xip
linux9s:~ # /sbin/xipinit-fw
Usage: init 0123456SsQqAaBbCcUu
linux9s:~ # cat /proc/mounts
rootfs / rootfs rw 0 0
/dev/root / ext3 rw 0 0
proc /proc proc rw 0 0
sysfs /sys sysfs rw 0 0
devpts /dev/pts devpts rw 0 0
tmpfs /dev/shm tmpfs rw 0 0
none /xip xip2 ro 0 0
none /lib/libdl.so.2 xip2 ro 0 0
none /lib/tls/librt.so.1 xip2 ro 0 0
none /lib/tls/libpthread.so.0 xip2 ro 0 0
none /lib/tls/libm.so.6 xip2 ro 0 0
none /lib/tls/libc.so.6 xip2 ro 0 0
none /lib/libz.so.1.2.1 xip2 ro 0 0
none /lib/ld-2.3.3.so xip2 ro 0 0
none /usr/local/sbin/snmpd xip2 ro 0 0
none /usr/local/lib/libnetsnmpmibs.so.5.2.1 xip2 ro 0 0
none /usr/local/lib/libnetsnmpagent.so.5.2.1 xip2 ro 0 0
none /usr/local/lib/libnetsnmphelpers.so.5.2.1 xip2 ro 0 0
none /usr/local/lib/libnetsnmp.so.5.2.1 xip2 ro 0 0
none /usr/lib/librpm-4.1.so xip2 ro 0 0
none /usr/lib/librpmdb-4.1.so xip2 ro 0 0
none /usr/lib/librpmio-4.1.so xip2 ro 0 0
none /usr/lib/libpopt.so.0.0.0 xip2 ro 0 0
none /usr/lib/libbz2.so.1.0.0 xip2 ro 0 0
```

Once we know that the overmount script works properly, we have to modify Linux so that this overmounting is done at boot time. We change the call to **init** in "zipl.conf" to use the overmount script "xipinit-fw" (see Example 2-32). This script will run at boot time to do the overmounting and then execute the standard Linux **init** routine.

Example 2-32 Update zipl.conf with the call to xipinit-fw and run zipl

```
linux9s:/etc #cd /etc
linux9s:/etc #vi zipl.conf
parameters = "dasd=0.0.0205,0.0.0306 root=/dev/dasda1 selinux=0
TERM=dumb elevator=cfq mem=256M init=/sbin/xipinit-fw"
linux9s:/etc # zipl -V
Using config file '/etc/zipl.conf'
Target device information
Device.....: 5e:00
Partition.....: 5e:01
Device name.....: dasda
DASD device number.....: 0205
Type.....: disk partition
Disk layout.....: ECKD/compatible disk layout
Geometry - heads.....: 15
Geometry - sectors.....: 12
Geometry - cylinders.....: 1113
Geometry - start.....: 24
File system block size.....: 4096
Physical block size.....: 4096
Device size in physical blocks...: 200316
Building bootmap in '/boot/zipl'
Adding IPL section 'ipl' (default)
kernel image.....: /boot/image at 0x10000
kernel parmline...: 'dasd=0.0.0205,0.0.0306 root=/dev/dasda1 selinux=0
TERM=dumb elevator=cfq mem=256M init=/sbin/xipinit-fw' at 0x1000
initial ramdisk...: /boot/initrd at 0x800000
Preparing boot device: dasda (0205).
Syncing disks...
Done.
```

This can be tested by rebooting the system with the **reboot** command.

You should see something like Example 2-33 on the virtual machine console during the boot of the system. This indicates that the DCSS has been loaded and the files are being mounted.

Example 2-33 Messages during boot

```
extmem info:segment_load: loaded segment NETSNMP range 0f800000 ..
0ffffff typ SR in shared mode
```

When the system is back up, check to see if the modules are mounted.
Example 2-34 demonstrates that the **mount** command does not show them:

Example 2-34 Mount command does not show overmounts

```
linux9s:~ # mount
/dev/dasda1 on / type ext3 (rw,acl,user_xattr)
proc on /proc type proc (rw)
sysfs on /sys type sysfs (rw)
tmpfs on /dev/shm type tmpfs (rw)
devpts on /dev/pts type devpts (rw,mode=0620,gid=5)
```

But using `cat /proc/mounts` will, as shown in Example 2-35:

Example 2-35 Use cat command to see overmounts

```
linux9s:~ # cat /proc/mounts
rootfs / rootfs rw 0 0
/dev/root / ext3 rw 0 0
none /xip xip2 ro 0 0
none /lib/libdl.so.2 xip2 ro 0 0
none /lib/tls/librt.so.1 xip2 ro 0 0
none /lib/tls/libpthread.so.0 xip2 ro 0 0
none /lib/tls/libm.so.6 xip2 ro 0 0
none /lib/tls/libc.so.6 xip2 ro 0 0
none /lib/libz.so.1.2.1 xip2 ro 0 0
none /lib/ld-2.3.3.so xip2 ro 0 0
none /usr/local/sbin/snmpd xip2 ro 0 0
none /usr/local/lib/libnetsnmpmibs.so.5.2.1 xip2 ro 0 0
none /usr/local/lib/libnetsnmpagent.so.5.2.1 xip2 ro 0 0
none /usr/local/lib/libnetsnmphelpers.so.5.2.1 xip2 ro 0 0
none /usr/local/lib/libnetsnmp.so.5.2.1 xip2 ro 0 0
none /usr/lib/librpm-4.1.so xip2 ro 0 0
none /usr/lib/librpmdb-4.1.so xip2 ro 0 0
none /usr/lib/librpmio-4.1.so xip2 ro 0 0
none /usr/lib/libpopt.so.0.0.0 xip2 ro 0 0
none /usr/lib/libbz2.so.1.0.0 xip2 ro 0 0
proc /proc proc rw 0 0
sysfs /sys sysfs rw 0 0
devpts /dev/pts devpts rw 0 0
tmpfs /dev/shm tmpfs rw 0 0
linux9s:~ #
```

Your system is now up and running with the SNMP daemon in an XIP2 DCSS. If you make the same changes to all your Linux images, then the snmp code will be shared among all these images.

Using DCSS for Oracle

This chapter discusses the method we used to set up a DCSS containing the Oracle RDBMS code, so that multiple Linux guests could share this code.

The method of implementing is very similar to the implementation of a DCSS as discussed in Chapter 2, however there are a few differences.

The documentation for the XIP2 file system is in *How to use Execute-in-Place Technology with Linux on z/VM*, SC33-8287, which you can find at:

http://www-128.ibm.com/developerworks/linux/linux390/october2005_documentation.html

This documentation, and the supporting sample shell scripts assume that the code to be loaded in the DCSS is located in the root file system. Because Oracle is an application, and is not in the root file system, we modified the sample shell scripts, and used a different method of initializing the DCSS. This chapter documents this approach.

3.1 Sizing the DCSS

After many experiments, we decided that the most effective way to use the DCSS is to share an entire directory, and only put the code that we knew would be executed into the DCSS. In our tests we obtained the best results by putting only two modules in the DCSS.

To make effective use of the DCSS, you should ensure that the DCSS is sized correctly. We decided to create a directory and put only the `oracle`, and `sqlplus` binaries in it. We used the `ls` command to determine the size of these modules, as demonstrated in Example 3-1.

Example 3-1 Get size of modules

```
pazxxt19:/local/10g/db/dcss # ls -al
total 87748
drwxr-xr-x  2 root  root    4096 Feb  9 17:39 .
drwxr-xr-x 62 oracle dba    4096 Dec  9 12:46 ..
-rwxr-x--x  1 oracle dba 89736931 Oct 28 09:26 oracle
-rwxr-x--x  1 oracle dba  12131 Oct 28 09:27 sqlplus
```

We decided that 100 MB was big enough to hold the two modules that we wanted to put into the DCSS.

3.2 Determining DCSS addresses

As the DCSS in Example 3-1 is 100 MB, the next step is to determine the starting address of the DCSS. For this example, the virtual machine size is 512 MB. To get the most benefit in this case from the DCSS, the virtual machine will be reduced in size by 100 MB and the DCSS will be placed just above the virtual machine. The DCSS starting address will be:

$$512 \text{ MB} - 100 \text{ MB} = 412 \text{ MB}$$

For creating the DCSS, the page number must be determined to supply to the `CP` command that creates the DCSS. The starting address of 412 MB must be converted to hex, and this then converted to the page number. The DCSS start address of 248 MB is address hex "F800000". The end address is hex "1FFFFFFF", or 512 MB minus 1. The page address is then the hexadecimal value truncated, so in this case, starting page is 19C00, and ending page is 1FFFF.

For different applications, a different sized DCSS could be more appropriate.

3.3 Tailoring Linux storage size to support the DCSS

The virtual machine should now be reduced by 100 MB, and then Linux needs to be configured to address the DCSS. The configuration file “/etc/zipl.conf” requires modification. To do this, we use the standard `vi`, as shown in Example 3-2.

Example 3-2 Changing the zipl.conf file

```
cd /etc
vi zipl.conf
```

Add “**mem=512M**” to the parameter line as shown in Example 3-3.

Example 3-3 zipl.conf file

```
parameters = "root=/dev/dasdb1 selinux=0 TERM=dumb elevator=deadline
dasd=200-209,300-30f,400-40f,fixedbuffers mem=512M"
```

Run `zipl` to update everything. You may want to use the `-V` parameter to get details on zipl's actions, as shown in Example 3-4.

Example 3-4 Output from zipl -V command

```
zipl -V
Using config file '/etc/zipl.conf'
Target device information
  Device.....: 5e:00
  Partition.....: 5e:01
  Device name.....: dasda
  DASD device number.....: 0200
  Type.....: disk partition
  Disk layout.....: ECKD/compatible disk layout
  Geometry - heads.....: 15
  Geometry - sectors.....: 12
  Geometry - cylinders.....: 200
  Geometry - start.....: 24
  File system block size.....: 4096
  Physical block size.....: 4096
  Device size in physical blocks..: 35976
Building bootmap in '/boot/zipl'
Adding IPL section 'ipl' (default)
  kernel image.....: /boot/image at 0x10000
  kernel parmline...: 'root=/dev/dasdb1 selinux=0 TERM=dumb elevator=deadline
dasd=200-209,300-30f,400-40f,fixedbuffers mem=512M' at 0x1000
  initial ramdisk...: /boot/initrd at 0x1000000
Preparing boot device: dasda (0200).
Syncing disks...
```

Done.

3.4 Creating the DCSS

We chose to create the DCSS from CMS. In order to have the authority to create and populate a DCSS, the z/VM directory entry for the CMS user must be defined as Class E privileges. There must also be a **NAMESAVE** statement in the directory entry. These privileges allow the CMS user to define the DCSS. In Example 3-5, we called the DCSS “ORADCSS”.

Example 3-5 Defining the DCSS

```
cp defseg oradcsc 19c00-1ffff sr
Saved segment ORADCSS was successfully defined in fileid 0501.
```

Validating the NSS is done with the **QUERY** command, or simply **Q**, as shown in Example 3-6.

Example 3-6 Saving the DCSS to create the storage image

```
#CP Q NSS MAP
FILE FILENAME FILETYPE MINSIZE BEGPAG ENDPAG TYPE CL #USERS PARMREGS VMGROUP
0401 LNXSEG1 DCSS N/A 54000 7FEFF SR A 00000 N/A N/A
0473 NETSNMP DCSS N/A 5C000 5FFFF SR A 00000 N/A N/A
...
0475 ORADCSS DCSS N/A 19C00 1FFFF SR S 00000 N/A N/A
```

Note that after the **saveseg** command in Example 3-7, the class of the DCSS has changed from “S” for Skeleton to “A” for Active.

Example 3-7 Results of the “saveseg oradcsc” command.

```
#CP saveseg oradcsc
Saved segment ORADCSS was successfully saved in fileid 0475.
CP Q NSS MAP
FILE FILENAME FILETYPE MINSIZE BEGPAG ENDPAG TYPE CL #USERS PARMREGS VMGROUP
0401 LNXSEG1 DCSS N/A 54000 7FEFF SR A 00000 N/A N/A
0473 NETSNMP DCSS N/A 5C000 5FFFF SR A 00000 N/A N/A
...
0475 ORADCSS DCSS N/A 60000 7FEFF SR A 00000 N/A N/A
```

3.5 Set the Linux size and boot Linux

Now that you have saved the DCSS, you have to change the size of your Linux virtual machine so that its pages don't overlap with the DCSS pages when the DCSS is loaded. To do this, either change the directory for the Linux virtual machine, or use the **define storage command in the Linux machine**. Changing the virtual machine size with the **define storage** command works immediately, but the change will be lost if you logoff the virtual machine.

After your system is booted, check to ensure the “mem=” parameter is present and correct by using **cmdline**, as shown in Example 3-8.

Example 3-8 Verify zipl changes were done correctly

```
pazxxt19:~ # cat /proc/cmdline
dasd=0.0.0205,0.0.0306 root=/dev/dasda1 selinux=0 TERM=dumb elevator=cfq
mem=512M BOOT_IMAGE=0
```

3.6 Resize your virtual machine and reboot

Now that you have saved the DCSS, you have to change the size of your virtual machine so that its pages will not overlap with the DCSS pages when the DCSS is loaded. To do this, shutdown your system, resize the virtual machine, and reboot. Issue the commands shown in Example 3-9 to do this Linux.

Example 3-9 Shutdown, change storage size and reboot

```
pazxxt19:~ # shutdown -h now
Broadcast message from root (pts/0) (Thu Dec 22 10:57:27 2005):
The system is going down for system halt NOW!
...
The system will be halted immediately.

...

Power down.
HCPGSP2630I The virtual machine is placed in CP mode due to a SIGP stop
and store status from CPU 00.
DEF STOR 412M
STORAGE = 0412M
Storage cleared - system reset.
IPL 205
```

Changing the virtual machine size with the **define storage** command works immediately, but the change will be lost if you logoff the virtual machine. To make this a permanent change, update the CP directory entry for this virtual machine.

The size of the machine is on the **USER** statement. This will ensure that the virtual machine size will be correct through system restarts.

After your system is back up, ensure the `mem=` parm is present and correct by using the `cmdline` command, as shown in Example 3-10.

Example 3-10 Verify zipl changes were done correctly

```
pazxxt19:~ # cat /proc/cmdline
dasd=0.0.0205,0.0.0306 root=/dev/dasda1 selinux=0 TERM=dumb elevator=cfq
mem=256M BOOT_IMAGE=0
```

3.7 Setup for the XIP2 DCSS

This DCSS differs from the SNMP DCSS discussed in 2.7, “Setup for the XIP2 DCSS” on page 12, in that it is being defined in an LVM outside of the root file system. We can use a similar technique to fill the DCSS with the programs to be shared, but we cannot use the same technique to initialize the DCSS after booting. In Example 3-11, we will create a directory to hold the modules we want to populate the DCSS.

We then create a directory under “oracle\home” called “dcss” in which we move the modules that we will be sharing with “dcss”. Before the move, we copy the modules for backup purposes. Once we move the modules we create a symbolic link to the new location of these modules.

Example 3-11 Create the directory for the mount point

```
mkdir $ORACLE_HOME/dcss
cp $ORACLE_HOME/bin/oracle $ORACLE_HOME/bin/oracle.org
cp $ORACLE_HOME/bin/sqlplus $ORACLE_HOME/bin/sqlplus.org
mv oracle $ORACLE_HOME/dcss/oracle
mv sqlplus $ORACLE_HOME/dcss/sqlplus
ln -s $ORACLE_HOME/dcss/oracle $ORACLE_HOME/bin/oracle
ln -s $ORACLE_HOME/dcss/sqlplus $ORACLE_HOME/bin/sqlplus
```

Now load the DCSS driver with the `modprobe` command. This driver is supplied in the Linux distribution, and allows you to treat the DCSS as a disk so that you can read and write data to it. We will use this driver to create a filesystem on the DCSS and copy the files we want to share into this filesystem. Example 3-12 shows the `modprobe` command that will load the block driver.

Example 3-12 Load the DCSS block mode driver

```
pazxxt19:~ # modprobe dcssblk
```

In order to define devices to use this driver, we need to define some nodes. To do this we need the driver's major number. To determine this number we look at the `/proc/devices` file, as shown in Example 3-13. We see that the `dcssblk` driver is there and the major number is 252.

Example 3-13 Find the driver's major number

```
pazxxt19:~ # cat /proc/devices
```

```
Character devices:
```

```
 1 mem
 4 ttyS
 4 ttysclp
 5 /dev/tty
 5 /dev/console
 5 /dev/ptmx
 9 st
10 misc
21 sg
128 ptm
136 pts
227 ttyTUB
228 fs3270
```

```
Block devices:
```

```
 1 ramdisk
 7 loop
 8 sd
 9 md
11 sr
65 sd
66 sd
67 sd
68 sd
69 sd
70 sd
71 sd
94 dasd
128 sd
129 sd
130 sd
131 sd
132 sd
133 sd
134 sd
135 sd
252 dcssblk
253 device-mapper
254 mdp
```

The next step is to create the device nodes, if this is required. Example 3-14 demonstrates the command used to verify whether or not the device nodes need to be created. In this example, we see that the output from the `ls -l /dev/dcss*` command returns “No such file or directory”, so the next step would be to create the device nodes, as indicated in the example by the use of the `mknod` commands for each node that we want to create.

Example 3-14 Verify device nodes and create if necessary

```
pazxxt19:~ # ls -l /dev/dcss*
/bin/ls: /dev/dcss*: No such file or directory

pazxxt19:~ # mknod /dev/dcssblk0 b 252 0
pazxxt19:~ # mknod /dev/dcssblk1 b 252 1
pazxxt19:~ # mknod /dev/dcssblk2 b 252 2
pazxxt19:~ # mknod /dev/dcssblk3 b 252 3

pazxxt19:~ # ls -l /dev/dcss*
brw-r--r-- 1 root root 252, 0 Oct 12 17:15 /dev/dcssblk0
brw-r--r-- 1 root root 252, 1 Oct 12 17:16 /dev/dcssblk1
brw-r--r-- 1 root root 252, 2 Oct 12 17:16 /dev/dcssblk2
brw-r--r-- 1 root root 252, 3 Oct 12 17:16 /dev/dcssblk3
```

3.8 Load the DCSS storage with the files to be shared

At this point, we have created a DCSS segment that will be loaded from address 412 MB to 512 MB. The machine is defined to be 412 MB in size and rebooted at this size. The Linux command line has been modified to support addresses up to 512 MB.

The `dcssblk` driver is used to manipulate the DCSS (Example 3-15).

Example 3-15 Use the `dcssblk` driver to load the DCSS into storage

```
extmem info:segment_load: loaded segment ORADCSS range 0000000019c00000 ..
000000001fffffff type SR in shared mode
dcssblk info: Loaded segment ORADCSS, size = 104857600 Byte, capacity = 204800
(512 Byte) sectors
Feb  9 17:37:02 pazxxt19 kernel: extmem info:segment_load: loaded segment
ORADCSS range 0000000019c00000 .. 000000001fffffff type SR in shared mode
Feb  9 17:37:02 pazxxt19 kernel: dcssblk info: Loaded segment ORADCSS, size =
104857600 Byte, capacity = 204800 (512 Byte)sectors
```

We now must get exclusive access to the DCSS while we are writing into it. We again use the `dcssblk` driver to do this (Example 3-16).

Example 3-16 Get exclusive access to the DCSS

```
pazxxt19:~ # echo 0 > /sys/devices/dcssblk/ORADCSS/shared
```

Now that we have exclusive access, we can define a filesystem on the DCSS (Example 3-17).

Example 3-17 Make a filesystem on the DCSS

```
pazxxt19:~ # mke2fs -b 4096 /dev/dcssblk0
mke2fs 1.36 (09-Feb-2005)
Filesystem label= OS type: Linux
Block size=4096 (log=2)
Fragment size=4096 (log=2)
2048 inodes, 2048 blocks
102 blocks (4.98%) reserved for the super user
First data block=0
1 block group
32768 blocks per group, 32768 fragments per group
2048 inodes per group
Writing inode tables: done
Writing superblocks and filesystem accounting information: done
This filesystem will be automatically checked every 37 mounts or
180 days, whichever comes first. Use tune2fs -c or -i to override.
pazxxt19:~ #
```

Example 3-18 shows how we mount the new filesystem under `/mnt` to create the directory structure and copy the modules. We need to make sure we have the same directory structure we have in the Oracle home directory, this is why we will be creating `local/10g/db` under `/mnt` so when we mount the DCSS filesystem it will overmount on `/local/10g/db/dcss`, where `/local/10g/db` is the Oracle home directory.

Example 3-18 Mount new filesystem

```
pazxxt19:~ # mount /dev/dcssblk0 /local/10g/db/dcss
pazxxt19:~ # cp -va /dcss/* /local/10g/db/dcss

pazxxt19:~# export ORACLE_HOME=/local/10g/db
pazxxt19:~ # mount /dev/dcssblk0 /mnt
pazxxt19:~ # cd /mnt
pazxxt19:/mnt # mkdir -p local/10g/db/dcss
pazxxt19:/mnt # cp -av $ORACLE_HOME/dcss/* local/10g/db/dcss
~/local/10g/db/dcss/oracle' -> ~/local/10g/db/dcss/oracle'
```

```
`/local/10g/db/dcss/sqlplus' ->
`local/10g/db/dcss/sqlplus'
```

Example 3-19 shows the next step, saving the DCSS. This will take the data in the current storage copy of the DCSS and schedule it for writing out to the z/VM SPOOL. Again we use the dcssblk driver to do this.

Example 3-19 Save the segment

```
pazxxt19:/ # echo 1 > /sys/devices/dcssblk/ORADCSS/save
dcssblk info: Segment ORADCSS is currently busy, it will be saved when it
become
s idle...
Feb 20 14:24:39 pazxxt19 kernel: dcssblk info: Segment ORADCSS is currently
busy , it will be saved when it becomes idle...
```

Unmount the file system as shown in Example 3-20. This means that no one is connected to the DCSS any more, so it will be written out.

Example 3-20 Free the segment so that it will be written

```
pazxxt19:/ # umount /mnt
HCPNSD440I Saved segment ORADCSS was successfully defined in fileid 0506.
HCPNSS440I Saved segment ORADCSS was successfully saved in fileid 0506.
dcssblk info: Segment ORADCSS became idle and is being saved now
Feb 20 14:26:33 pazxxt19 kernel: dcssblk info: Segment ORADCSS became idle and
is being saved now
```

We have now finished using the segment as a disk. We can release it, and then mount it using the new XIP2 filesystem (Example 3-21).

Example 3-21 Free the DCSS after populating it

```
echo ORADCSS> /sys/devices/dcssblk/remove
pazxxt19:/dcss # ls -l /sys/devices/dcss*
total 0
drwxr-xr-x 2 root root 0 Dec 12 17:25 .
drwxr-xr-x 7 root root 0 Dec 12 16:56 ..
--w----- 1 root root 0 Dec 12 17:06 add
-rw-r--r-- 1 root root 4096 Dec 12 17:06 detach_state
--w----- 1 root root 0 Dec 12 17:25 remove
```

Unlike the example shown in Chapter 2, “Putting the SNMP daemon into a DCSS” on page 5, we are sharing a complete directory in this DCSS. The example script “xipinit.sh” is the supplied code which initializes a DCSS, doing the “overmounting” that makes the files available. We copy this file to /etc/init.d

with a new name, and modify it to point to our environment. We then modify rc2.d, rc3.d, and rc5.d to cause these scripts to be executed after the boot sequence (see Example 3-22).

Example 3-22 Update xipinit.sh and put it into /etc/init.d

```
pazxxt19:~ # cd /etc/init.d
pazxxt19:/etc/init.d # cp /local/scripts/xip-howto-scripts/xipinit.sh
init.oradcsc
vi init.dcss
```

using vi set:

```
#mount point of xipimage
ROMOUNT="/oraxip"
#name of xipimage
XIPIIMAGE="ORADCSS"
#name of device node
RODEV="/dev/dcssb1k0"

RODIRS=/local/10g/db/dcss
# make sure it ends with ,
RODIRS="$RODIRS",
```

Also inset a # in the last line to make the call to init a comment:
run real init
#\$FAKE exec /sbin/init "\$@"

In directory /etc/init.d, there are directories rc2.d, rc3.d and rc5.d which contain a sequence of scripts that get executed when Linux enters runlevel 2, 3 or 5. When we installed Oracle on this Linux machine, a file named "S17init.cssd" was created in these three directories. This shell script initializes the cssd component of Oracle. Our shell, "script init.oradcsc", must get invoked before the cssd.

initialization. These initialization scripts get executed in sequence. So we rename the Oracle script from “S17init.cssd” to “S18init.cssd” and insert our script as “S17init.oradcss” (see Example 3-23).

Example 3-23 Rename the init call for Oracle cssd

```
pazxxt19:/etc/init.d # find . -name S17init.cssd
./rc2.d/S17init.cssd
./rc3.d/S17init.cssd
./rc5.d/S17init.cssd

pazxxt19:/etc/init.d # mv ./rc2.d/S17init.cssd ./rc2.d/S18init.cssd
pazxxt19:/etc/init.d # mv ./rc3.d/S17init.cssd ./rc3.d/S18init.cssd
pazxxt19:/etc/init.d # mv ./rc5.d/S17init.cssd ./rc5.d/S18init.cssd
```

Example 3-24 shows how we insert a call to our initialization routine as “S17init.oradcss”. This will get executed before the initialization of Oracle cssd because it is now named “S18init.cssd”.

Example 3-24 Insert call to our initialization routine in sequence

```
pazxxt19:/etc/init.d # ln -s /etc/init.d/init.oradcss ./rc2.d/S17init.oradcss
pazxxt19:/etc/init.d # ln -s /etc/init.d/init.oradcss ./rc3.d/S17init.oradcss
pazxxt19:/etc/init.d # ln -s /etc/init.d/init.oradcss ./rc5.d/S17init.oradcss
```

A reboot of the machine should load and initialize the Oracle DCSS. Example 3-25 demonstrates this.

Example 3-25 Boot up message

```
<notice>start services (local init.oradcss)
Feb 20 14:44:13 pazxxt19 kernel: extmem info:segment_load: loaded segment
ORADCSS range 0000000019c00000 .. 000000001fffffff type SR in shared mode.
```

3.9 Use same segment in a “hole” in memory

Now that this segment is built to load from 412 MB to 512 MB, it could be used by other, larger Linux guests. We had to add mem=512M to the “zipl.conf” file so that Linux could address up to 512 MB even though the machine was only 412 MB in size. If we remove the mem= parameter from “zipl.conf”, Linux will be able to address the size of the machine where it is IPLed.

z/VM allows you to define storage with a “hole” in the physical memory using the config form of the define storage command as shown in Example 3-26.

Example 3-26 Defining storage with a hole

```
def stor config 0.412M 512M.1536m
STORAGE = 1948M
Storage Configuration:
0.412M 512M.1536M
Extent Specification          Address Range
-----
          0.412M             0000000000000000 - 0000000019BFFFFFF
          512M.1536M         0000000020000000 - 000000007FFFFFFF
```

A machine defined this way can load the same ORADCSS as created earlier in this chapter because the DCSS will be loaded in the hole in the address space. The mem= parameter on the Linux guest should not be specified. With the above definition of storage Linux will be able to address up to 2 GB.

Archived

DCSS measurements

This chapter describes the tests and measurements we made to show that DCSS was actually saving memory with multiple guests when using SNMP and Oracle.

4.1 Oracle and XIP2 objectives

The execute in place (XIP2) filesystem allows Linux servers to share programs, executing the code directly out of the DCSS. The objective is that with many Linux servers sharing storage, programs only need to be resident once and should allow for significant storage savings.

There are trade-offs that must be understood before implementing XIP2. The trade-off is between the costs of addressing additional storage, the cost of the DCSS, and then the savings in Linux storage. At a minimum, the savings in Linux storage *must* be larger than the costs in addressing the DCSS. Otherwise, storage will be lost for each server utilizing the DCSS, missing our objective. The improvement in Linux storage for one server will not exceed the cost of the DCSS itself, but will be hopefully covered on the second or third server utilizing the DCSS. If the savings are so low and the cost of the DCSS is very high, then it may take more servers utilizing the DCSS than exist, again missing the objective.

Recognizing that the objectives are:

- ▶ Minimizing the cost of storage addressability
- ▶ Minimizing the cost of the DCSS
- ▶ Maximizing the Linux storage savings

Each of these should be measured, managed and are not completely intuitive. Much of the analysis is to understand Linux storage, since the first measurements using “common knowledge” proved to not reduce storage as much as expected. It was found that performance improvements using XIP2 could not be taken for granted.

ESALPS Version 4.1, the Linux Performance Suite was provided by Velocity Software for this analysis. More information about ESALPS can be found in the Appendix A, “ESALPS” on page 61.

4.2 Experiment objectives

Installations running a few large (2 GB - 6 GB) servers will not see a large gain from implementing an XIP2 file system in a DCSS. The objective of our experimentation, therefore, is for installations that want to run 20 or 30 Oracle servers in the 256M to 512M range will see the real benefit from adding servers without adding more real storage.

4.3 XIP2 analysis findings

For Oracle servers, the XIP2 filesystem in a DCSS was shown to reduce storage when created correctly. It was also found that the implementation could create storage overheads that exceeded the benefit. Most of the study was in determining how to optimize the XIP2 implementation. This is the first time this kind of analysis has been published, possibly the first time ever done. There were several surprises that, without measurements, would never have been realized. It became obvious that measuring both the Linux storage down to the process level and the z/VM storage in a constrained environment gave a different set of answers than had been expected.

There are two ways that the XIP2 feature may be used. One method is to load an entire directory into the DCSS. This was found to be very ineffective due both to the high cost of addressability and the inflexibility of what files can be shared.

The other is to load individual files into the DCSS. This was found to be most effective in using the XIP2 DCSS feature. The DCSS should be loaded with only frequently used programs and it should have a minimal amount of free space. This is slightly more complex, but provides much more measurable benefits.

Another finding was in where to place the DCSS. Three methods of placing a DCSS may be used:

- ▶ Adding a DCSS at the address just above the Linux servers
- ▶ Adding a DCSS at the highest possible storage
- ▶ Putting the DCSS into a “hole”. Current z/VM allows defining a virtual machine with a gap in the middle of virtual machines and thus reserving a place for the DCSS. For installations that want to implement a DCSS when using machines that are 2 GB or larger, this is the only option because, at least up to z/VM V5.2, DCSS must be below the 2 GB line.

Each method has reasons for doing it that way. However, adding a DCSS to an existing Linux server adds storage to the Linux server and no savings will be obtained, in fact the storage requirements will increase. Linux will always use all available storage, the only way to force a reduction in storage requirements is to reduce the virtual machine storage size. In implementing a DCSS with XIP2, the Linux server storage **MUST** be reduced, with a recommendation of reducing the storage by the size of the DCSS.

The implementation most likely to have minimal storage savings is when adding the DCSS in high storage.

The most significant benefit found in these experiments using XIP2 and DCSS was to carefully pack the DCSS with only frequently used programs, place the

DCSS in as low a storage position as possible, and then reduce the size of the virtual machine.

The value in including libraries in the DCSS is likely positive, but further experiments are needed to validate this benefit.

4.4 Linux storage

This section examines storage overheads and how to measure them.

4.4.1 XIP2 storage overheads

There are three measurable costs associated with using a DCSS for XIP2 that can be greater than the benefits without management. Two of the three measurable costs are related to addressability. To address additional storage, Linux needs to increase its page management structure tables and page tables. Thus, having a DCSS in high storage may cost more storage in overhead than save in shared storage.

4.4.2 Storage addressability overhead

In the XIP2 in DCSS measurements, the “mem=” parameter was used to increase the amount of storage that Linux can address when adding a DCSS outside the virtual machine. This parameter builds the page management structures required to address additional storage that Linux perceives when it boots. This would not be necessary if using the “hole” approach. Several experiments were done with just changing the “mem=” parameter to show the storage costs in a static environment and to understand the storage costs.

- ▶ Addressability overhead shows up in the size of the kernel. This increases due to page management structure tables. Measurements showed the following costs for 31-bit Linux:
- ▶ Page management structures (struct page entries) for 31-bit Linux overhead is 8 MB per gigabyte of addressed storage (measured in kernel storage). The entries are 32 bytes per page.
- ▶ An extra megabyte per gigabyte of overhead was identified, but the cause of this was not determined.

The costs using 64-bit Linux are double that of when using 31-bit Linux and therefore needs to be managed even more carefully. Page management structure overhead is 16 MB per gigabyte (measured in kernel storage). The entries are 64 bytes per page.

4.4.3 Linux storage analysis

ESALPS provides several measures of storage. In the following example of an ESAUCD2 report, the measures of storage are:

- ▶ Kernel - this includes the kernel and page management structures. This is the difference in size between what Linux reports as “total Real Storage” and the virtual machine size.
- ▶ Available - storage that is not in use by Linux. This is always low. The only value that this number has is when it is large; this shows that Linux has not been able to find any data to put into storage.
- ▶ Buffer - storage in the write buffer, typically small
- ▶ Cache - dynamic storage in use for caching data and other Linux storage items. Note that this includes the Oracle SGA, as well as all executables.
- ▶ Overhead - storage in use by control blocks and page tables.

An experiment was performed in analyzing just the impact of “mem=”. The following report shows two 31-bit SUSELNX servers that are identical except for the defined memory size. Note the difference in the total real storage size, by increasing addressability to 1G from 256M; the total storage size has decreased 8 MB, the kernel size has increased 8 MB.

Example 4-1 ESAUCD2 report

Report: ESAUCD2	LINUX EXAMPLE			Linux Test				
Node/	-----Mega Bytes-----							
Time/	<--Real Storage-->			. <----Storage in Use----->				
Date	Total	Avail	Used	. Shared	Buffer	Cache	Ovrhd	
10:26:00				.				
SUSELNX1	239.3	175.8	63.5	.	0	18.6	26.1	18.8
SUSELNX2	247.3	182.8	64.5	.	0	18.5	26.2	19.8

The kernel size is the difference between the virtual machine size, in this case, 256 MB for both servers and the reported Total Real Storage size.

What this analysis shows is that if a DCSS was defined at the 512M mark for 512M, the mem= would have to be 1.G, adding a cost of storage of 8 MB to an existing 256 MB 31-bit server.

4.5 Linux virtual storage analysis

Before analyzing storage impacts of the DCSS, it was noticed that the mem= parameter in Linux had impacted Linux storage requirements in different areas. The kernel size increased for one, as did the overhead storage. The following table shows the results of several simple experiments measuring the impact of “mem=”.

These measurements were taken to help understand how to measure the impact of increasing Linux addressability. ESALPS was used to show the kernel size and the overhead storage size for several different configurations first in 31-bit mode, and then in 64-bit mode.

All values in the following tables are in megabytes.

4.5.1 31-bit Linux analysis

Table 4-1 Memory analysis

mem=	Kernel Size	Cache Size	Buffer Size	Overhead Storage
256	8.3	25.1	3.8	16.7
256	8.3	25.2	3.7	16.8
512	10.6	18.9	3.9	18.0
528	10.8	25.2	3.7	19.6
1024	15.2	25.2	3.8	19.6
1040	15.3	25.2	4.0	22.2
1552	19.8	25.1	4.1	22.4

In this study, the virtual machine was 256 MB. The objective was to determine addressability costs. By analyzing this data, and the slope of the line, the kernel size increases at a rate of 9 MB per GB of defined storage. The overhead storage also increased at a rate of 4 MB per GB. Thus, adding a 256 MB DCSS to a 256 MB virtual machine would cost about 3.5 MB of Linux virtual storage. It was later determined that this storage is referenced and not paged out by VM.

4.5.2 64-bit Linux analysis

From the following two samples extracted from an ESAUCD2 report, the first set is three machines, all with machine size of 1024, and mem=2047. The second set is 512 MB servers, but with mem=1024. The static storage then for the kernel

is 43 MB for the first, and 27 MB for the second. Arithmetic says that the cost of the 1 GB extra of memory addressability costs 16 MB of Linux virtual storage, and further arithmetic would indicate that the fixed cost of the 64-bit kernel is about 11 MB (27 MB - 16 MB).

Example 4-2 ESAUCD2 report

```
Report: ESAUCD2      LINUX U
Monitor initialized:
-----
Node/    <-----
Time/    <--Real Storage-->
Date     Total Avail Used
-----
PAZXXT19 981.1 826.3 154.8
PAZXXT21 981.1 751.1 230.0
PAZXXT20 981.1 776.4 204.7
-----
PAZXXT19 485.1  10.3 474.8
PAZXXT20 485.1  11.6 473.5
PAZXXT21 485.1  10.2 475.0
```

4.6 VM real storage analysis

One potential error in the mem= analysis is that if Linux does not reference pages, then the extra page management structure tables are paged out by z/VM and would not take storage other than during initialization.

4.6.1 Measuring mem= impact on VM working set

The following analysis shows first with 64 MB addressable, and the second with 1024 MB. Both measurements were done with a 64 MB Linux server, with the server being tested running a minimal configuration, including netnmp for instrumentation. A second server was started with a very large virtual machine that caused memory contention to force our test server to be paged out, leaving only the referenced pages resident.

The test then came when a new process was started to see if there was an impact on storage that was different between the two tests. The "top" process was used as its characteristics should be very familiar to Linux users.

This analysis shows that the VM working set increases about 8 MB when the mem=1024M is used, as expected for 31-bit Linux. 64-bit Linux would be double.

Therefore using a DCSS at an address just below the 2 GB line for a 512 MB server for a 64-bit server will add 24 MB to its virtual storage requirements *and* its working set. Thus the savings must be much more than 24 MB per server for this to be worthwhile. If the same DCSS was compressed to 100 MB, and put at the 512 line, then the added overhead would be less than 2 MB, and much more significant savings are then possible.

4.6.2 Test 1, mem=64M

The following report shows the active process analysis during the experiment. The ESALNXP report shows each active process, its process ID, parent ID, CPU and storage information as provided in the Linux process table. Following is the z/VM storage analysis. Note that the SNMP daemon and some background Linux tasks (pdflush, slpd) were running throughout the measurement. Then top was started at 14:57. This made bash and sshd active (to run the terminal session), touching a lot of pages.

The interval at 14:55 has the least amount of processes running and should have the lowest working set.

Example 4-3 ESALNXP report for test 1

Report: ESALNXP LINUX HOST Process Statistics Report

node/ Name	<-Process Ident->			<----CPU Percents----->						<Stg (k)>	
	ID	PPID	GRP	Tot	sys	user	syst	usrt	nice	Size	RSS

14:53:00											
SUSELNX1	0	0	0	9.5	8.6	1.0	0	0	0	135K	53K
pdflush	1509	4	0	0.1	0.1	0	0	0	0	0	0
snmpd	1817	1	1816	2.6	2.3	0.3	0	0	0	7060	4016
slpd	1832	1	1832	0.1	0.1	0.0	0	0	0	3588	1244

14:54:00											
SUSELNX1	0	0	0	10.0	8.1	1.8	0	0	0	135K	53K
init	1	0	0	0.5	0.4	0.0	0	0	0	628	100
kblockd/	6	4	0	0.1	0.1	0	0	0	0	0	0
pdflush	1509	4	0	0.4	0.4	0	0	0	0	0	0
snmpd	1817	1	1816	4.0	3.0	1.0	0	0	0	7060	4016
slpd	1832	1	1832	0.2	0.2	0.0	0	0	0	3588	1244

14:55:00											
SUSELNX1	0	0	0	6.3	4.0	2.3	0	0	0	135K	53K
snmpd	1817	1	1816	2.6	1.6	1.0	0	0	0	7060	4016

14:56:00											
SUSELNX1	0	0	0	9.7	6.2	3.5	0	0	0	135K	53K


```

snmpd      1817      1 1816 3.8 2.4 1.4  0  0  0 7060 4016
sshd       2073 1868 2073 0.3 0.3 0.0  0  0  0 8392 2576

```

14:57:00

```

SUSELNX1      0      0      0 13.3 8.8 4.5  0  0  0 137K 54K
kjournal     277      1      1 0.2 0.2  0  0  0  0  0  0
snmpd        1817      1 1816 2.4 1.7 0.8  0  0  0  0 7060 4016
sshd         2073 1868 2073 0.6 0.4 0.2  0  0  0 8392 2580
bash         2076 2073 2076 0.6 0.5 0.1  0  0  0 3204 1952
top          2095 2076 2095 2.0 1.1 0.9  0  0  0 2132 1100

```

14:58:00

```

SUSELNX1      0      0      0 19.9 12.5 7.4  0  0  0 137K 54K
kjournal     277      1      1 0.1 0.1  0  0  0  0  0  0
snmpd        1817      1 1816 5.2 3.8 1.3  0  0  0 7060 4016
sshd         2073 1868 2073 0.7 0.4 0.3  0  0  0 8392 2580
top          2095 2076 2095 2.7 1.1 1.6  0  0  0 2132 1100

```

The next part of the test was to show what happens on the VM side. Note that the numbers used for Linux are in “k”, and the numbers for z/VM are in pages (4k).

The following ESAUSR2 report shows resource requirements by virtual machine. In analyzing the SUSELNX1 server during the experiment, its total resident storage started at 16K pages, the full 64 MB of the virtual machine. When there was contention this dropped to 1160 pages which would be just the referenced pages for when just Linux and the SNMP daemon were running. 1160 pages is assumed to be the minimum working set for this configuration.

When pdflush started, the VM resident pages went up to 1460 pages, and when top started, the resident pages went up to 4242 pages. When top starts, top runs under bash, and sshd runs to validate the logon.

Example 4-4 ESAUSR2 report for test 1

```

Report: ESAUSR2      User Resource      Linux Test
-----
      <---CPU time--> <-Pages--> <-----Paging (pages)----->
UserID <(seconds)> T:V <Resident> <---Allocated---> <---I/O--->
/Class  Total  Virt Rat Totl Activ  Total ExStg  Disk  Read Write
-----
14:47:00
SUSELNX1 2.48 2.03 1.2 16K 16222      0  0  0  0  0
(All storage resident)
14:48:00
SUSELNX1 2.47 2.03 1.2 16K 16222      0  0  0  0  0
14:49:00
SUSELNX1 3.05 2.28 1.3 5584 5584 11037  907 10130 1021 10145

```

```

(Contention starts, pushes 10,130 pages to disk)
14:50:00
SUSELNX1  4.98  4.15  1.2  1879  1879  14937    21 14916  2916  5033
14:51:00
SUSELNX1  2.68  1.92  1.4  1189  1189  15658    197 15461  2974  2028
14:52:00
SUSELNX1  5.63  5.01  1.1  1196  1196  15754    366 15388  3082  2271
14:53:00
SUSELNX1  3.89  3.13  1.2  1160  1160  15819    410 15409  2431  1907
(Minimum storage requirement, working set starts to stabilize)
14:54:00
SUSELNX1  3.33  2.63  1.3  1461  1461  15498    195 15303    143    52
(kernel processes start: init,kblockd,pdflush)
14:55:00
SUSELNX1  3.33  2.67  1.2  1331  1331  15630    362 15268    37     0
(storage starts to drop)
14:56:00
SUSELNX1  3.70  2.94  1.3  3910  3910  15405    144 15261  2361    0
(Start "top" - cost 10MB)
14:57:00
SUSELNX1  5.04  4.40  1.1  4135  4135  15056    136 14920    217    0

```

The conclusion from this is that our minimal configuration running just Linux and netsnmp without a DCSS was about 1160 pages. When top was started, from the z/VM perspective, the working set adds about 3000 pages. Linux reports the combined size of top, bash and sshd as about 5.6 MB or 1400 pages. The remainder of the working set increase is not completely explained but likely in large part due to kernel data areas being referenced by top.

When all processes were active and part of the working set, the working set was 4135 pages; this is what we hope to reduce using XIP2.

4.6.3 Test 2, mem=1024M

In the second test with mem=1024M to validate the numbers, the increase when top was started was consistent. The storage requirements in this data is consistently 8 MB greater. There is no difference in extra costs when top was started. The base cost is about 3000 pages instead of the 1160 pages prior, or about 7.6 MB. This is consistent with the increase in kernel size due to the page management structure tables.

This validates the conclusion that the cost in addressing additional storage impacts both real storage and virtual storage.

Example 4-5 ESALNXP report for Test2

```

Report: ESALNXP      LINUX HOST Process Statistics Report

```

```

-----
node/      <-Process Ident-> <-----CPU Percents-----> <Stg (k)>
Name      ID   PPID  GRP  Tot  sys user syst usrt nice  Size RSS
-----
16:30:00
SUSELNX1  0    0    0   5.3  3.4  1.9  0  0  0 198K  78K
snmpd    1823  1  1822  1.7  1.1  0.6  0  0  0 7060 4016
-----
16:31:00
SUSELNX1  0    0    0   6.8  4.8  1.9  0  0  0 198K  78K
snmpd    1823  1  1822  2.2  1.6  0.6  0  0  0 7060 4016
-----
16:32:00
SUSELNX1  0    0    0   6.3  3.9  2.4  0  0  0 198K  78K
snmpd    1823  1  1822  2.1  1.3  0.8  0  0  0 7060 4016
-----
16:33:00
SUSELNX1  0    0    0   5.1  3.3  1.8  0  0  0 198K  78K
snmpd    1823  1  1822  1.6  1.0  0.6  0  0  0 7060 4016
-----
16:36:00
SUSELNX1  0    0    0  14.4  8.6  5.8  0  0  0 200K  79K
snmpd    1823  1  1822  1.7  1.1  0.6  0  0  0 7060 4016
sshd     2082 1875 2082  0.7  0.5  0.2  0  0  0 8392 2580
bash     2085 2082 2085  0.3  0.3  0.0  0  0  0 3204 1952
top      2104 2085 2104  2.1  1.0  1.1  0  0  0 2132 1100
-----
16:37:00
SUSELNX1  0    0    0  13.1  7.6  5.5  0  0  0 200K  79K
snmpd    1823  1  1822  1.5  0.9  0.6  0  0  0 7060 4016
sshd     2082 1875 2082  0.8  0.6  0.1  0  0  0 8392 2580
top      2104 2085 2104  2.0  0.9  1.2  0  0  0 2132 1100
-----

```

The following user report shows the resident storage that decreased to referenced pages only, and then increased when top was started.

Example 4-6 ESAUSR2 report for test2

```

-----
Report: ESAUSR2      User Resource U      Linux Test
Monitor initialized: First record analyzed: 01/
-----
      <---CPU time--> <-----Ma <-----Paging (pages)----->
UserID <(seconds)> T:V <Resident> <---Allocated---> <---I/O--->
/Class  Total  Virt Rat Totl Activ  Total ExStg  Disk  Read Write
-----
16:25:00
SUSELNX1 2.56 2.09 1.2 16K 15876      0  0  0  0  0
16:26:00

```

SUSELNX1	2.51	2.08	1.2	10K	10430	5446	5371	75	0	75
16:27:00										
SUSELNX1	3.33	2.86	1.2	3725	3725	12481	9137	3344	330	3269
16:28:00										
SUSELNX1	2.29	1.85	1.2	2984	2984	13345	9135	4210	123	866
(working set at lowest point with just netsnmp)										
16:29:00										
SUSELNX1	2.27	1.84	1.2	2997	2997	13339	9129	4210	65	0
16:30:00										
SUSELNX1	2.26	1.83	1.2	3109	3109	13339	9129	4210	54	0
16:31:00										
SUSELNX1	2.40	1.94	1.2	3169	3169	13339	9129	4210	60	0
16:32:00										
SUSELNX1	2.28	1.85	1.2	3169	3169	13339	9129	4210	0	0
16:33:00										
SUSELNX1	2.26	1.83	1.2	3169	3169	13339	9129	4210	0	0
(working set is very stable with just netsnmp)										
16:34:00										
SUSELNX1	2.29	1.85	1.2	3745	3745	13327	9117	4210	564	0
16:35:00										
SUSELNX1	2.35	1.89	1.2	5354	5354	13391	9181	4210	1673	0=
(top is started, working set increases about 3200 pages)										
16:36:00										
SUSELNX1	4.02	3.44	1.2	6011	6011	13840	9630	4210	1106	0
16:37:00										
SUSELNX1	3.96	3.41	1.2	6022	6022	13829	9619	4210	0	0
16:38:00										
SUSELNX1	3.93	3.39	1.2	6027	6027	13824	9614	4210	0	0

4.6.4 Implementing XIP2 in DCSS to reduce Linux Resident Storage

The next step was to look for the “best case” opportunity to minimize the overall real storage requirements by implementing XIP2 in a DCSS.

The methodology used for packing the DCSS was to evaluate the ESALNXC report for this server to determine which processes should be loaded into a DCSS. This report shows all the resident processes, not just the active ones, and their parent/child relationship. Child processes are indented from the parent. The RSS, Resident Storage, is the relevant number. This is the amount of virtual storage currently addressable by the process. Of the resident storage, some will be private, and some “shared”, which can then be in the DCSS. Only experimentation and measurements will show the savings. Of these processes, the idle processes do not impact the referenced pages that VM will keep resident. Therefore, putting inactive processes into a DCSS has little benefit unless used by multiple servers. Note that processes with a zero RSS are kernel processes and would not be included in a DCSS.

Example 4-7 ESALNXC report

```

Report: ESALNXC          LINUX Process Conf Report
-----
Node/      <-Process Ident-> <Process> <Storage(k)
Name       ID    PPID  GRP  Path      Size  RSS
-----
SUSELNX1
init        1     0     0  init [3]   628   99
*           4     1     0                0     0
khelper     5     4     0  khelper    0     0
kblockd/    6     4     0  kblockd/   0     0
pdflush    1516  4     0                0     0
kjournal    277   1     1  pdflush    0     0
rc          557   1     557 /sbin/sy   2616  144
S14hwsca   1921  557   557 top         2616  256
  hwbootsc 1929  1921  557 sshd: ro    2612  208
  hwscan    1953  1929  557 -bash      2880  1064
blogd      568   1     568 /sbin/kl   9824   76
syslogd    1798  1     1798 /sbin/kl   1640   728
snmpd      1817  1     1816 /sbin/re   7060  4016
resmgrd    1830  1     1830 /sbin/po   1496   516
portmap    1831  1     1831 /usr/sbi   1516   580
syslogd    1841  1     1841 /sbin/re   1640   196
klogd      1844  1     1844 /sbin/po   1596   224
sshd       1868  1     1868 /sbin/mi   5304  1972
  sshd     2073  1868  2073 -bash      8392  2580
  bash     2076  2073  2076 top         3204  1952
  top     2095  2076  2095 sshd: ro    2132  1104
  sshd    6524  1868  6524 -bash      8392  2576
  bash    6527  6524  6527 pdflush    3208  1996
-----

```

From this report, four processes were chosen for inclusion into a DCSS, snmpd, top, bash and sshd. These four were used during the previous experiments, setting a baseline and should provide a measurable impact. The target storage was the approximately 4200 pages when both snmpd and top (and sshd) were running.

The four processes required a 13 MB DCSS to be created, which was located at 64 MB, just above the Linux virtual machine. The following sample when all four processes were running shows that each process was greatly reduced in resident virtual storage.

Example 4-8 ESALNXP report for test 2

```

Report: ESALNXP          LINUX HOST Process Statistics Report
-----
node/      <-Process Ident-> <-----CPU Percents-----> <Stg (k)>
Name       ID    PPID  GRP  Tot  sys user syst usrt nice  Size RSS
-----

```

SUSELNX1	0	0	0	10.6	6.2	4.4	0	0	0	131K	15K
kjournal	279	1	1	0.1	0.1	0	0	0	0	0	0
snmpd	1865	1	1864	1.3	0.7	0.6	0	0	0	7248	1948
sshd	2125	1923	2125	0.6	0.5	0.2	0	0	0	8392	740
bash	2128	2125	2128	0.2	0.2	0.0	0	0	0	3204	592
top	3171	2128	3171	3.3	1.7	1.5	0	0	0	2132	288

In measuring the results, from the Linux perspective, there was a savings in process resident virtual storage of about 4.7 MB.

Table 4-2 Results

Process	Without DCSS	With DCSS	Savings (MB)
snmpd	4016K	1948K	2.1
sshd	2580	740K	.5
bash	1952	592K	1.3
top	1104	288K	.8
Total:			4.7

The snmpd process went from 4016K storage to 1948K storage, a savings of 2.1 MB. Thus with a 13 MB dcss, the virtual storage savings was 4.7 MB.

4.6.5 Impact of DCSS on z/VM real storage

The first analysis was with just the SNMP daemon running. The pages resident for z/VM varied much more in this experiment than prior experiments. We believe that the use of XIP2 increased the use of the pdflush kernel process, increasing the storage requirements, with storage varying between 500 and 1400 pages. The overall storage requirement in the equivalent experiment without XIP2 varied from 1160 to 1460 pages. Thus the savings was as much as 600 pages, 2.4 MB. The conclusion from this experiment is that the savings in the SNMP daemon virtual storage closely matched the savings in the working set.

Example 4-9 ESUSR2 report

Report: ESAUSR2		User Resource Linux Test									
		<---CPU time--> <-Pages-->					<-----Paging (pages)----->				
UserID	<(seconds)>	T:V	<Resident>	<---Allocated---	<---I/O---	Total	ExStg	Disk	Read	Write	
/Class	Total	Virt	Rat	Totl	Activ	Total	ExStg	Disk	Read	Write	
SUSELNX1	4.21	3.63	1.2	16K	16186	0	0	0	0	0	
16:29:00											
SUSELNX1	5.04	4.46	1.1	16K	16186	0	0	0	0	0	

```

16:30:00
SUSELNX1 2.90 2.31 1.3 1290 1290 15046 12465 2581 641 2759
16:31:00
SUSELNX1 5.59 5.04 1.1 1198 1198 15657 4912 10745 3548 10838
16:32:00
SUSELNX1 3.55 2.89 1.2 785 785 16417 1051 15366 2675 6475
16:33:00
SUSELNX1 5.90 5.35 1.1 1111 1111 16548 1206 15342 2547 1813
16:34:00
SUSELNX1 3.26 2.56 1.3 981 981 16667 1342 15325 35 15
16:35:00
SUSELNX1 4.64 3.96 1.2 1402 1402 16505 1232 15273 311 0
16:36:00
SUSELNX1 3.37 2.69 1.3 925 925 17015 1709 15306 89 89
16:37:00
SUSELNX1 4.68 3.99 1.2 738 738 17373 1993 15380 795 678
16:38:00
SUSELNX1 4.63 3.95 1.2 498 498 17639 2342 15297 155 48

```

Note the different processes running at different intervals. As these are more or less normal processes, their overhead is to be expected.

Example 4-10 ESALNXP report

Report: ESALNXP LINUX HOST Process Statistics Report

node/ Name	<-Process ID	Ident-> PPID	GRP	<----CPU Percents----->						<Stg (k)>	
				Tot	sys	user	syst	usrt	nice	Size	RSS
SUSELNX1	0	0	0	5.4	4.4	0.9	0	0	0	129K	15K
kjournal	279	1	1	0.1	0.1	0	0	0	0	0	0
snmpd	1865	1	1864	4.0	3.2	0.8	0	0	0	7248	1948
slpd	1887	1	1887	0.1	0.1	0	0	0	0	3588	436

16:33:00											
SUSELNX1	0	0	0	9.0	4.5	0.8	1.9	0.7	0	129K	15K
init	1	0	0	0.2	0.2	0	0	0	0	628	116
khelper	5	4	0	3.0	0.4	0	1.9	0.7	0	0	0
pdflush	99	4	0	0.2	0.2	0	0	0	0	0	0
kjournal	279	1	1	0.3	0.3	0	0	0	0	0	0
snmpd	1865	1	1864	4.0	3.1	0.8	0	0	0	7248	1948
slpd	1887	1	1887	0.3	0.3	0.0	0	0	0	3588	436

The relevant information from this test was just saving the SNMP daemon in a DCSS could be shown to reduce real storage requirements by as much as 2 MB.

The second part of the analysis was to show the value of the other three processes. The top process was started at 16:44. Note the jump in storage to

2550 pages, but after initialization this storage drops down and then varies between 1321 and 2559 pages. The equivalent prior measurement had a working set of between 3900-4200 pages. The saving represented from these four processes then is about 2500-2600 pages, or about 10 MB, from a 13 MB dcsc. With the variability in how Linux kernel processes run, this number is likely extremely optimistic as the virtual storage savings are much smaller than the real storage savings. But it is very obvious that the storage savings from the DCSS was significant in terms of percent of working set.

Example 4-11 ESAUSR2 report

```

Report: ESAUSR2      User Resource U      Linux Test
Monitor initialized:      First record analyzed: 01/
-----
      <---CPU time--> <-----Ma <-----Paging (pages)----->
UserID <(seconds)> T:V <Resident> <---Allocated---> <---I/O--->
/Class  Total  Virt Rat Totl Activ  Total ExStg  Disk  Read Write
-----
16:37:00
SUSELNX1 4.68  3.99 1.2  738  738 17373  1993 15380  795  678
16:38:00
SUSELNX1 4.63  3.95 1.2  498  498 17639  2342 15297  155  48
16:39:00
SUSELNX1 4.42  3.76 1.2 1119 1119 17527  2283 15244    1  0
16:44:00
SUSELNX1 5.82  5.03 1.2 2550 2550 18346  3163 15183 2311  0
16:45:00
SUSELNX1 5.04  4.27 1.2 1559 1559 19321  4283 15038  129  0
16:46:00
SUSELNX1 6.33  5.56 1.1 1321 1321 19461  4492 14969   39  0
16:47:00
SUSELNX1 6.31  5.52 1.1 1321 1321 19482  4528 14954   69  33
16:48:00
SUSELNX1 4.79  4.09 1.2 1420 1420 19407  4462 14945   33  0
-----

```

4.6.6 XIP2 DCSS guidelines from non-Oracle test

The following guidelines were derived from this test:

- ▶ Evaluate the processes to include in the DCSS from the ESALNXP and ESALNXC reports; picking the proper candidates will result in reducing storage requirements.
- ▶ Pick the highly used processes. The SNMP daemon is always a good candidate when ESALPS is used.
- ▶ Carefully pack the DCSS with these processes to minimize the size of DCSS.

- ▶ Locate the DCSS in as low a storage position as possible to reduce overhead.
- ▶ Measure it. *Note*, in several early experiments, no differences in storage were perceived. As the DCSS had been implemented “by the book”, the first conclusion was that XIP2 had no value. Resulting analysis showed that there were errors in the implementation scripts and the DCSS and XIP2 were in fact not being used. Only the measurements caught this problem.

4.7 Packing the XIP2 DCSS for Oracle

The first experiments were with small servers, and now the following is Oracle in larger servers.

4.7.1 Linux storage management

Linux storage management is designed to be efficient. When a program is loaded, only the pages referenced are made resident in virtual storage. If two programs load the same storage, then the pages used by both processes are shared and do not consume extra storage as more processes utilize the already loaded pages.

Thus for Oracle, in the following (abbreviated) process report, there are many Oracle processes, but with current instrumentation levels of Linux, it is unclear to how much storage is actually shared, and how much is resident in total to support Oracle. From this data that comes from the standard Linux process table, all we know is how much storage a given process can address currently. In looking at the Oracle processes, the largest in this case has 87M resident (87K * 1024). The other processes are sharing some amount of storage, of which potentially most of is part of the 87M.

Example 4-12 ESLNXP report

```
Report: ESALNXP          LINUX HOST Process Statistics Report
Monitor initialized: 01/12/06 at 15:27:11 on 2084
```

node/ Name	<-Process Ident->			<-----CPU Percents----->						<Stg (k)>	
	ID	PPID	GRP	Tot	sys	user	syst	usrt	nice	Size	RSS
PAZXXT21	0	0	0	0.4	0.3	0.0	0	0	0	12M	1M
snmpd	2571	1	2570	0.3	0.3	0.0	0	0	0	9784	5644
oracle	7848	1	7848	0	0	0	0	0	0	282K	15K
oracle	7850	1	7850	0	0	0	0	0	0	281K	13K
oracle	7852	1	7852	0	0	0	0	0	0	284K	33K
oracle	7854	1	7854	0.0	0.0	0	0	0	0	296K	15K
oracle	7856	1	7856	0	0	0	0	0	0	282K	17K

oracle	7858	1	7858	0	0	0	0	0	0	283K	52K
oracle	7860	1	7860	0	0	0	0	0	0	281K	20K
oracle	7862	1	7862	0	0	0	0	0	0	282K	31K
oracle	7864	1	7864	0	0	0	0	0	0	282K	13K
oracle	7866	1	7866	0	0	0	0	0	0	282K	12K
oracle	7872	1	7872	0	0	0	0	0	0	281K	14K
oracle	7876	1	7876	0	0	0	0	0	0	285K	58K
oracle	7878	1	7878	0	0	0	0	0	0	281K	24K
oracle	7880	1	7880	0	0	0	0	0	0	282K	87K

There are two methods of determining what files should be moved to the XIP2 DCSS. One method which was used for the Oracle measurements was to move a complete file directory containing all of the object code for Oracle. The DCSS was created at the 1.5 GB address for 512 MB. It appeared that storage requirements actually increased in using XIP2 in a DCSS. Now we know that the extra addressability cost at least 24 MB, and that the overall savings in shared storage was not much more than that.

The second and better method is to select the programs used most by the Linux servers and include those programs into a DCSS.

Note that any installation using the scripts provided on the IBM Web site may have encountered a bug, in which case users could have believed they were using the programs in shared mode from the DCSS, but were in fact using the programs off disk instead. This has been corrected on the IBM download page as a result of this research. Version 2 of the XIPINIT-fw shell script should be used.

4.7.2 Impact of DCSS on z/VM real storage

The pages in the DCSS that are referenced will stay resident. Pages in the DCSS that are in use by more than one Linux server result in savings. Pages that are rarely referenced may increase the storage requirements with little savings. The number of pages of the DCSS resident as well as the number of Linux servers with the DCSS attached (and in use) is provided on the ESADCSS display and report. In this case, the ORADCSS is a 512 MB DCSS (131,000 pages) with 26,000 pages resident.

Note that if this was a constrained environment, some amount of the DCSS would have been paged out.

Example 4-13 ESADCSS report

```
Report: ESADCSS      NSS/DCSS Analysis Linux
-----
                                <-Users-> <-----Number of Pag
```

Name	Spool ID	<Creation-> Date	<Creation-> Time	<First Use> Date	<First Use> Time	Non-Shrd	Non-Shrd	Saved	Non---resident-Data <2GB	>2GB	HO
13:34:00											
CMS	469	08/24	09:09	10/30	09:06	33	0	1298	0	687	0
CMSFILES	454	05/23	11:12	10/30	09:06	4	0	768	0	301	0
CMSPIPES	451	05/23	11:12	10/30	09:06	36	0	256	0	94	0
CMSVMLIB	458	08/22	15:43	10/30	09:06	36	0	256	0	113	0
ESAMON	348	11/10	13:21	10/30	09:06	4	0	0	16K	9772	0
GCS	440	05/20	13:33	10/30	09:06	3	0	120	917	939	0
HELPSSEG	466	08/22	16:12	10/31	13:42	1	0	256	0	230	0
INSTSEG	455	05/23	11:11	10/30	09:06	36	0	768	0	298	0
MONDCSS	350	11/10	13:32	10/30	09:06	2	0	0	16K	106	0
NLSAMENG	464	08/22	15:44	10/30	09:06	25	0	256	0	2	0
ORADCSS	492	11/07	13:04	11/07	13:11	3	0	131K	0	26K	0
SCEE	446	05/23	11:12	10/30	09:06	6	0	256	0	27	0
SCEEX	439	05/23	11:12	10/30	09:06	6	0	1792	0	279	0
SVM	453	05/23	11:12	10/30	09:06	2	0	256	0	39	0
VTAM	149	04/17	11:06	10/30	09:06	3	0	256	0	70	0
Totals:						200	0	138K	34K	39K	0

4.8 Measuring XIP2 - file directories

The following experiments were done using 512 MB virtual machines, a 512 MB DCSS, and a 220 MB SGA. The DCSS was created using a file directory approach rather than selecting specific programs, and at the 512 MB address.

The kernel size change is very measurable, and very consistent. The overhead storage measurement is dependant on the point in time of the script where the measurement was taken. There is a cost of 9 MB for the added 512 MB addressability.

The two ESALNXP report extracts show just the largest Oracle process in terms of resident storage to show before and after the DCSS was implemented. The first report extract shows that the resident storage grows over time and peaks at 184 MB. In the next experiment, it peaked at 146 MB, a 40 MB improvement.

Example 4-14 ESALNXP report

Report: ESALNXP LINUX HOST Process Statistics Report										
node/	<-Process Ident->			<-----CPU Percents----->					<Stg (k)>	
Name	ID	PPID	GRP	Tot	sys	user	syst	usrt	Size	RSS
oracle	3358	1	3358	0.2	0.2	0.0	0	0	370K	46K

```

oracle 3358 1 3358 0.4 0.3 0.1 0 0 370K 63K
oracle 3358 1 3358 0.8 0.4 0.4 0 0 370K 101K
oracle 3358 1 3358 0.8 0.4 0.4 0 0 370K 141K
oracle 3358 1 3358 0.9 0.4 0.5 0 0 370K 176K
oracle 3358 1 3358 0.9 0.4 0.5 0 0 370K 183K
oracle 3358 1 3358 1.0 0.4 0.5 0 0 370K 184K
oracle 3358 1 3358 0.9 0.4 0.5 0 0 370K 184K
oracle 3358 1 3358 0.8 0.3 0.5 0 0 370K 184K
oracle 3358 1 3358 0.8 0.4 0.4 0 0 370K 184K

```

Report: ESALNXP LINUX HOST Process Statistics Report
Monitor initialized: 01/12/06 at 15:27:11 on 2084

```

-----
node/      <-Process Ident-> <-----CPU Percents-----> <Stg (k)>
Name      ID   PPID  GRP  Tot  sys user syst usrt  Size RSS
-----
oracle    3158  1   3158  0.8  0.5  0.3  0   0 370K  59K
oracle    3158  1   3158  1.2  0.7  0.5  0   0 370K  101K
oracle    3158  1   3158  0.4  0.2  0.2  0   0 370K  117K
oracle    3158  1   3158  0.2  0.1  0.1  0   0 370K  146K

```

The results of this experiment show that savings were not as large as had hoped and actually non-existent for this data. Addressing the DCSS cost 9 MB per server, that we know is real storage, the cost of the 512 MB DCSS was measured to have 30 MB resident. And as we did not reduce the servers virtual storage, Linux immediately consumed any storage that was now available and used it for caching data. The lessons learned from this measurement are:

- ▶ Manage the contents of the DCSS. Creating large DCSS without contents that are highly used or unreferenced costs in Kernel storage without benefit.
- ▶ Minimizing the size of the DCSS and packing it with highly shared files will provide the most benefit.
- ▶ Place the DCSS in as low a storage address as possible.
- ▶ Reduce the virtual machine size. For example, if a 100 MB DCSS was used, then drop the virtual machine size from 512 MB to 412 MB, with the DCSS loaded at address 412 MB.

The conclusion then is that best case for Oracle, with a 100 MB DCSS, the cost is about 1 MB for addressability, plus the resident pages of the DCSS. Realistically, if we can load the DCSS with Oracle, SQL as well as the SNMP daemon and other processes utilized, the virtual storage requirements might drop 60 MB in which case reducing the virtual machine size by 60 MB would be reasonable. For a 512 MB server, this is over a 10% savings, and in a constrained environment where z/VM is removing unreferenced pages is likely much more.

Note that none of the Oracle measurements were in a storage constrained environment so that there is no valid comparisons of the VM resident storage, this should be in a follow on project.

4.9 Measuring XIP2 - direct files

This experiment has two servers. the T20 server has an active DCSS, the T21 server has no DCSS. All servers are 400 MB. The DCSS is 100 MB and starts at 400 MB. The DCSS was packed with just two components, the Oracle and sqlplus processes.

For this experiment there was no contention, so from a VM perspective there is no difference in resident storage requirements. In all cases, the T20 and T21 machines had all 400M (100K pages) resident.

Example 4-15 ESAUSR2 report

```

Report: ESAUSR2      User Resource Utilization Linux Test
-----
      <---CPU time--> <-----Main St <-----Paging (pages)----->
UserID <(seconds)> T:V <Resident> Lock <---Allocated---> <---I/O--->
/Class  Total  Virt Rat Totl Activ -ed Total ExStg Disk Read Write
-----
14:42:00 145.3 143.9 1.0 5.0M 5013K 1190    20M 2605K 17M  929 2705
PAZXXT20 15.02 14.96 1.0 101K 101K 14.0     34  34   0   0   0
PAZXXT21 13.26 13.15 1.0 101K 101K 13.0     5   5   0   0   0
-----
14:43:00 279.6 278.1 1.0 5.0M 5014K 1189    20M 2605K 17M 1144   0
PAZXXT20 58.73 58.62 1.0 101K 101K 13.0     19  19   0   0   0
PAZXXT21 58.69 58.57 1.0 101K 101K 13.0     2   2   0   0   0
-----
14:44:00 276.0 274.6 1.0 5.0M 5014K 1189    20M 2604K 17M  656  600
PAZXXT20 57.91 57.82 1.0 101K 101K 13.0     21  21   0   0   0
PAZXXT21 56.95 56.83 1.0 101K 101K 13.0     2   2   0   0   0
-----
14:45:00 202.5 201.2 1.0 5.0M 5011K 1189    20M 2604K 17M 1089 2255
PAZXXT21 37.61 37.49 1.0 101K 101K 13.0     3   3   0   0   0
PAZXXT20 33.59 33.53 1.0 101K 101K 13.0    20  20   0   0   0
-----

```

One sample of the process data report looks like the following. With the Oracle processes sharing storage, it is difficult to come up with extremely accurate conclusions. However, these three servers ran exactly the same workload. If you take the largest Oracle process (in the RSS/resident storage) column, the T21 server is 60M and the T20 server is 30M, which would make a 30 MB reduction in virtual storage requirements seem likely.

Example 4-16 ESALNXP Report

Report: ESALNXP LINUX HOST Process Statistics Report											
node/ Name	<-Process Ident->			<-----CPU Percents----->						<Stg (k)>	
	ID	PPID	GRP	Tot	sys	user	syst	usrt	nice	Size	RSS
14:42:00											
PAZXXT21	0	0	0	5.1	0.6	0.1	1.2	3.2	0	6M	682K
init	1	0	0	4.3	0	0	1.1	3.2	0	696	288
kswapd0	232	1	1	0.2	0.2	0	0	0	0	0	0
snmpd	2571	1	2570	0.4	0.3	0.1	0	0	0	9784	5644
nscd	2859	1	2859	0.0	0.0	0	0	0	0	43K	912
bash	3258	3257	3258	0.0	0.0	0	0.0	0	0	4788	2360
oracle	3288	1	3288	0.1	0.1	0.0	0	0	0	296K	15K
oracle	3290	1	3290	0.0	0.0	0	0	0	0	281K	15K
oracle	3313	1	3313	0.0	0	0.0	0	0	0	283K	60K
PAZXXT20	0	0	0	12.4	0.4	0.2	2.6	9.3	0	15M	1M
init	1	0	0	11.9	0.0	0	2.6	9.3	0	696	288
pdflush	231	4	0	0.0	0.0	0	0	0	0	0	0
kswapd0	232	1	1	0.0	0.0	0	0	0	0	0	0
snmpd	2575	1	2574	0.2	0.2	0.0	0	0	0	9788	5648
bash	3265	3264	3265	0.0	0	0.0	0	0.0	0	4788	2360
oracle	3295	1	3295	0.2	0.1	0.1	0	0	0	296K	7088
oracle	3320	1	3320	0.0	0	0.0	0	0	0	282K	30K

To look at several points over time, the following report shows only the Oracle and sqlplus processes. For the T21 server, the Oracle (process ID 3313) process starts at 60M and goes to 67M. For the T20 server, the Oracle (process ID 3320) process go from 30M to 38M. As this is the same workload, a savings again of 30M of virtual storage appears obvious. The sqlplus is a much smaller process, it added another 1.6M in savings.

Example 4-17 ESALNXP report

Report: ESALNXP LINUX HOST Process Statistics Report											
node/ Name	<-Process Ident->			<-----CPU Percents----->						<Stg (k)>	
	ID	PPID	GRP	Tot	sys	user	syst	usrt	nice	Size	RSS
PAZXXT21	0	0	0	5.1	0.6	0.1	1.2	3.2	0	6M	682K
oracle	3313	1	3313	0.0	0	0.0	0	0	0	283K	60K
PAZXXT20	0	0	0	12.4	0.4	0.2	2.6	9.3	0	15M	1M
oracle	3320	1	3320	0.0	0	0.0	0	0	0	282K	30K
14:43:00											
PAZXXT21	0	0	0	99.2	1.7	0.9	27.2	69.4	0	6M	720K
oracle	3286	1	3286	0.1	0.1	0.1	0	0	0	284K	38K
PAZXXT20	0	0	0	99.4	1.4	0.4	26.5	71.1	0	16M	1M
oracle	3316	1	3316	0.0	0	0.0	0	0	0	285K	35K

14:44:00											
PAZXXT21	0	0	0	100	1.2	0.6	18.9	79.7	0	7M	769K
oracle	3313	1	3313	0.2	0.0	0.2	0	0	0	283K	65K
sqlplus	6869	3328	3325	0.0	0	0.0	0	0	0	21K	6196
PAZXXT20	0	0	0	98.6	1.0	0.8	18.0	78.8	0	16M	1M
oracle	3320	1	3320	0.3	0.0	0.3	0	0	0	282K	36K
sqlplus	6976	3325	3324	0.0	0.0	0.0	0	0	0	22K	7840
14:45:00											
PAZXXT21	0	0	0	1.9	1.0	0.9	0	0	0	7M	777K
oracle	3313	1	3313	0.2	0.0	0.2	0	0	0	283K	67K
PAZXXT20	0	0	0	1.5	0.7	0.8	0	0	0	16M	1M
oracle	3320	1	3320	0.2	0	0.2	0	0	0	282K	38K

The cost of the DCSS for this was 18K pages. As this was NOT a storage constrained environment, the cost is likely not this high.

Example 4-18 ESADCSS report

Report: ESADCSS NSS/DCSS Analysis											
Name	Spool ID	Creation		First Use		Users		Non-		Number of	
		Date	Time	Date	Time	Shrd	Shrd	Saved	Data	<2GB	>2GB

14:42:00											
CMS	469	08/24	09:09	01/12	14:38	33	0	1298	0	491	0
CMSFILES	454	05/23	11:12	01/12	14:38	4	0	768	0	144	0
CMSPIPES	451	05/23	11:12	01/12	14:38	36	0	256	0	79	0
CMSVMLIB	458	08/22	15:43	01/12	14:38	36	0	256	0	67	0
ESAMON	348	11/10	13:21	01/12	14:38	4	0	0	16K	2293	0
GCS	440	05/20	13:33	01/12	14:38	4	0	120	917	613	0
INSTSEG	455	05/23	11:11	01/12	14:38	36	0	768	0	122	0
MONDCSS	350	11/10	13:32	01/12	14:38	2	0	0	16K	145	0
NLSAMENG	464	08/22	15:44	01/12	14:38	25	0	256	0	2	0
ORADCSS	500	01/11	12:19	01/12	14:38	2	0	25600	0	18K	0
SCEE	446	05/23	11:12	01/12	14:38	6	0	256	0	14	0
SCEEX	439	05/23	11:12	01/12	14:38	6	0	1792	0	35	0
SVM	453	05/23	11:12	01/12	14:38	3	0	256	0	12	0
VTAM	149	04/17	11:06	01/12	14:38	4	0	256	0	66	0
Totals:						201	0	31882	34K	22K	0

The conclusion from this analysis duplicates prior conclusions that it is likely that at least a 10% in virtual storage can be attained, and though the real storage was not evaluated, it should be expected to reduce real storage by much more than that.

4.10 Conclusion

The savings are measurable and predictable. Further experiments should include analysis in a storage constrained environment. Also, including the operational libraries in the DCSS may also increase savings beyond what was measured.



ESALPS

For the measurements and analysis for this redbook, ESALPS (Linux Performance Suite) from Velocity Software was used for its unique capabilities and focuses on the z/VM and Linux environments.

ESALPS overview

ESALPS, the Linux Performance Suite, is a suite of products provided by Velocity Software for measurement and tuning of z/VM, Linux under z/VM, and applications in this environment. The four products that make up this suite are:

- ▶ ESAMAP, the VM Monitor Analysis Program, providing performance reports on all aspects of VM/ESA® and z/VM performance
- ▶ ESAMON, the VM Real Time Monitor, providing real time data collection and analysis of performance
- ▶ ESATCP, the network and Linux data collection program
- ▶ ESAWEB, a very fast VM based Web server

ESALPS addresses the infrastructure needs of the following:

- ▶ Performance: By providing full current and historical data for z/VM, Linux, and network performance characteristics, current performance and prior performance problems are easily analyzed and resolved. Reporting based on selecting specific periods of time allows in-depth performance analysis of performance problems.
- ▶ Capacity Planning: Long term data in the form of a performance database (PDB™) is needed as input to long term capacity planning and trend analysis. ESALPS meets the needs of all capacity planners by providing the PDB, as well as many interfaces to their favorite enterprise capacity management tools.
- ▶ Operations: ESALPS provides alerts on Linux, z/VM and network performance characteristics, either via 3270 interface, browser interface or via SNMP interface to the installation's management console.
- ▶ Accounting: By accurately accounting for CPU, storage and I/O, ESALPS can be used to account for virtual machines, or for Linux down to the process and application level.

ESALPS features

ESALPS provides the following features:

- ▶ Full function z/VM performance monitor allows for monitoring of all z/VM subsystems, including full reporting, real time analysis and management reporting.
- ▶ Network Monitor: Allows monitoring of your internet servers via SNMP.
- ▶ Linux Monitor: Allows monitoring of your Linux disks, storage and processes. In a virtual environment, ESALPS has the very unique capability for correcting

the process data that in a virtual environment are often incorrect, sometimes by an order of magnitude.

- ▶ Monitoring NT, SUN, HP: ESALPS will monitor any server that supports the standard MIBs. This includes network information, disk and storage information, and process data.
- ▶ Web serving: ESALPS includes a full function Web server to provide both browser access to the data but also to allow Web site support and CGI development on the VM platform.
- ▶ Alerts: Different types of alerts can be setup for network availability monitoring, VM subsystem errors and threshold exceptions, Linux errors and threshold exceptions.
- ▶ Full Performance Database: ESALPS provides both real time data and historical data for in-depth analysis. The performance data is collected daily with a one minute granularity based on the monitor interval. A longer term archive is collected usually with a granularity of 15 minutes. This PDB includes VM data, Linux data and network data.

Critical agent technology

One of the powerful features of ESALPS is the agent technology. Agents running inside Linux may be either active agents or passive agents. An active agent will wake up every few seconds on a regular schedule, collect data, possibly write data to a log, and then sleep. A passive agent will wait for an external data request and will not otherwise wake up and use resource.

The issue is that under z/VM, we expect to have maybe hundreds of Linux servers, of which many will be idle. With an active agent in each of them, the Linux servers are not idle and may consume excessive resources just by processing the active agent requests. Thus the only way to conserve resources is to eliminate the active agent from most or all of the servers.

With the passive technology enabled by using SNMP agents, ESALPS will know that a server is idle because of the z/VM data, and will not have to wake up an agent inside Linux, thus causing additional paging and CPU consumption, to determine and record the fact that the server is idle. The idle servers will not be measured until they become active.

The operational cost of agents may be quite high. On other non-virtualized servers, the cost is relatively insignificant as it only impacts the applications on that one server. In a virtualized environment where sharing resources is a primary objective, resources used by each agent reduces the resources available to all of the other servers. Minimizing the cost of running these agents

is critical to the success of monitoring many Linux servers under z/VM. The SNMP passive agent technology is believed to be the lowest cost agent technology available for Linux on z/VM.

NETSNMP extensions

The design of NETSNMP and SNMP host-specific MIBs was found lacking in a few areas. Velocity Software has extended the MIBs supported by NETSNMP. This agent can be downloaded from Velocity Software's Web site. This corrects measured values such as swap rate, and provides added process data thought necessary to meet the needs of this environment. This agent also is modified to use less resource than the NETSNMP that would come with the normal Linux distributions.

Standard interface

ESALPS uses standard interfaces for all data collection. The advantage to using the standard interfaces provided is that when there are a multitude of releases and distributions available, the standard interfaces will provide consistent data sources.

z/VM provides a "monitor interface" that has been available since 1988. Since then, this interface has provided a consistent view of performance of VM systems.

Network performance is collected using SNMP, the standard for network management.

NETSNMP, an open source software package provides host data for the Linux and other platforms.

VM Application data interface is used by applications to insert data into the monitor stream consistent with the standard monitor interface. ESATCP uses this interface to ensure consistent data collection that allows full integration of Linux and VM data.

Scripts used to set up DCSS

In this appendix we provide listings of the shell scripts we downloaded from developerWorks.

xipinit.sh

Example: B-1 xipinit.sh - Sets up to use a DCSS for a complete directory

```
#!/bin/sh
#
# xipinit.sh, Version 2
#
# (C) Copyright IBM Corp. 2002,2005
#
# /sbin/xipinit: use read only files from another file system
#
# default options

# internals
#set -v
#FAKE=echo

#mount point of xipimage
ROMOUNT=""
#name of xipimage
XIPIIMAGE=""
#name of device node
RODEV=""

RODIRS="/lib,/usr/lib,/usr/X11R6/lib,/bin,/sbin,/usr/X11R6/bin,/usr/bin,"
# make sure it ends with ,
RODIRS="$RODIRS",

mount -t sysfs none /sys
if [ ! -e /sys/devices/dcssblk ]; then
    echo "xipinit: loading dcssblk module"
    /sbin/modprobe dcssblk
fi
echo $XIPIIMAGE > /sys/devices/dcssblk/add

# mount ro file system to its mount point
echo "xipinit: mounting read-only segment"
$FAKE mount -t ext2 -o ro,xip "$RODEV" "$ROMOUNT"
# bind mount all ro dirs into rw filesystem
while [ -n "$RODIRS" ] ; do
    dir="${RODIRS%%,*}"
    RODIRS="${RODIRS#*,}"
    test -d "$dir" || continue
    echo "xipinit: binding directory" $dir
    $FAKE mount --bind "$ROMOUNT/$dir" "$dir"
done
umount /sys
```

```
# run real init
$FAKE exec /sbin/init "$@"
```

xipinit-fw.sh

Example: B-2 xipinit-fw - set up to use DCSS for selected files

```
#!/bin/bash
#
# xipinit-fw.sh, Version 3
#
# (C) Copyright IBM Corp. 2005

#mount point of xipimage
ROMOUNT=""
#name of xipimage
XIPIIMAGE=""
#name of device node
RODEV=""

mount -t sysfs none /sys
if [ ! -e /sys/devices/dcssblk ]; then
    echo "xipinit-fw: loading dcssblk module"
    /sbin/modprobe dcssblk
fi
echo $XIPIIMAGE > /sys/devices/dcssblk/add

echo "xipinit-fw: mounting read-only segment"
/bin/mount -t ext2 -o ro,xip $RODEV $ROMOUNT
echo "xipinit-fw: binding files"
for i in $(/usr/bin/find $ROMOUNT/*)
do
    if [[ -f $i ]]; then
        /bin/mount --bind $i ${i#$ROMOUNT}
    fi
done
umount /sys

exec /sbin/init $@
```

copylibs.sh

Example: B-3 *copylibs.sh*

```
#!/bin/bash
# copylibs.sh, Version 3
#
# (C) Copyright IBM Corp. 2005
#
#
# copies the binary and the libraries used to destination directory
# usage: ./copylibs.sh -f <executable|file> -d <destination directory>
#
#FAKE=echo

function getlibs {
    for i in `ldd $FILE | awk '{print $3}'`
    do
        echo $i
        if [[ -h $i ]]
        then
            echo $i is a link
            #LINKPATH=${i%/*}
            FILELINKEDTO=$(readlink -f $i)
            #FILELINKEDTOPATH=${FILELINKEDTO%/*}
            echo $FILELINKEDTO
            $FAKE cp -a --parent ${FILELINKEDTO} ${DESTINATION}
        elif [[ -e $i ]]
        then
            $FAKE cp -a --parent $i ${DESTINATION}
        fi
    done

    echo NOTE: Libraries which were loaded with libdl will not be copied.
    echo NOTE: be sure that you copy these extra.
}

###
###
# here the script starts
###
###

if [[ $# -ne 4 ]]
then
    echo "Usage: ./copylibs.sh -f <executable|file> -d <destination directory>"
    exit 1
fi
```



```

while getopts :f::d: OPT
do
  case $OPT in
    f )  if [[ ! $FILE && ${OPTARG:0:1} != '-' ]]
        then
            echo "option f is set with $OPTARG"
            FILE=$OPTARG
            echo $FILE
        else
            echo error for option -f please check
            exit 1
        fi
        ;;
    d )  if [[ ! $DESTINATION && ${OPTARG:0:1} != '-' ]]
        then
            echo "option d is set with $OPTARG"
            DESTINATION=$OPTARG
        echo $DESTINATION
        else
            echo error for option -d please check
            exit 1
        fi
        ;;
    :)   echo "no option set" ;exit 1 ;;
  esac
done

#here check if file (full qualified) exists
if [[ -e $FILE ]]
then
  if [[ $(file $FILE|grep ELF > /dev/null) -eq 0 ]]
  then
    getlibs
    # a dynamic shared object, so get all other libs needed
  fi
  #copy file itself first check if file is a symlink
  if [[ -h $FILE ]]
  then
    FILE=$(readlink -f $FILE)
  fi
  $FAKE cp -a --parent $FILE $DESTINATION
elif [[ ! -e $FILE ]]
then
  #check if file is an executable maybe found in PATH
  FILE=$(which $FILE)
  if [[ $? -ne 0 ]]
  then

```

```
        echo File not found, exiting...
        exit 1
    else
        $FAKE cp -a --parent $FILE $DESTINATION
        getlibs
    fi
fi
exit
```

Related publications

The publications listed in this section are considered particularly suitable for a more detailed discussion of the topics covered in this redbook.

IBM Redbooks

For information on ordering these publications, see “How to get IBM Redbooks” on page 71. Note that some of the documents referenced here may be available in softcopy only.

- ▶ *Experiences with Oracle Database on Linux on zSeries*, SG24-6552 (Oracle9i)
- ▶ *Linux for IBM zSeries and S/390: Distributions*, SG24-6264
- ▶ *Linux for S/390*, SG24-4987
- ▶ *Linux on IBM zSeries and S/390: ISP/ASP Solutions*, SG24-6299
- ▶ *Building Linux Systems under IBM VM*, REDP-0120
- ▶ *Linux on zSeries and S/390: Systems Management*, SG24-6820
- ▶ *z/VM and Linux on zSeries: From LPAR to Virtual Servers in Two Days*, SG24-6695
- ▶ *Linux Handbook A Guide to IBM Linux Solutions and Resources*, SG24-7000

Other publications

These publications are also relevant as further information sources:

- ▶ *How to use Execute-in-Place Technology with Linux on z/VM*, SC33-8287

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